

Tropical methods for ergodic control and zero-sum games

Minilecture, Part III

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Dynamical Optimization in PDE and Geometry
Applications to Hamilton-Jacobi
Ergodic Optimization, Weak KAM
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Today

Spectral theory

Algorithmic aspects

The max-plus spectral problem

Given $A = (A_{ij}) \in (\mathbb{R} \cup \{-\infty\})^{n \times n}$, find $v \in \mathbb{R} \cup \{-\infty\}^n$, $v \not\equiv -\infty$, $\lambda \in \mathbb{R}$, such that

$$\max_j A_{ij} + v_j = \lambda + v_i$$

$$\text{“}Av = \lambda v\text{”}$$

Among the oldest max-plus results.

Goes back to Cuninghame-Green 61, Vorobyev, Romanovski, Gondran and Minoux 77, Cohen, Dubois, Quadrat 83, ... Some references in Akian, SG, Bapat: Handbook of linear algebra (finite dim) and Max-plus Martin boundary / discrete spectral theory (infinite dim).

Interpretation: dynamic programming, one player

Set of nodes $[d] := \{1, \dots, d\}$, arc (i, j) with weight A_{ij}

$$A_{ij}^k = \sum_{m_1, \dots, m_{k-1} \in [d]} A_{im_1} A_{m_1 m_2} \cdots A_{m_{k-1} j}$$

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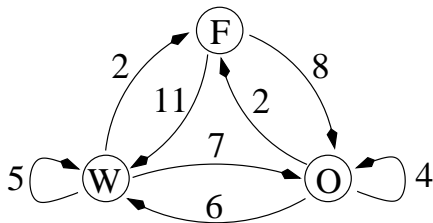
$$A_{ij}^k = \max_{m_1, \dots, m_{k-1} \in [d]} A_{im_1} + A_{m_1 m_2} + \dots + A_{m_{k-1} j}$$

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$$\begin{aligned} A_{ij}^k &= \max_{m_1, \dots, m_{k-1} \in [d]} A_{im_1} + A_{m_1 m_2} + \dots + A_{m_{k-1} j} \\ &= \max \text{ weight path } i \rightarrow j \text{ length } k \end{aligned}$$

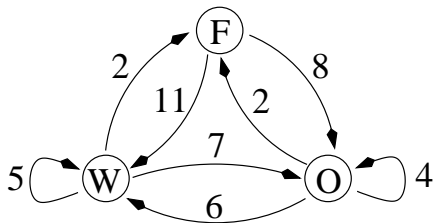
Crop rotation



A_{ij} = reward of the year if crop j follows crop i
F=fallow (no crop), W=wheat, O=oat,

$$(A^k v)_i = \sum_{j \in [d]} A_{ij}^k v_j$$

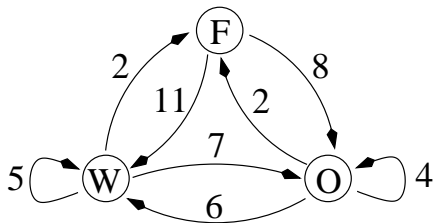
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= reward in k years, init. crop i ; v_j term. reward

Eigenvector

Find $v \in \mathbb{R}_{\max}^d$, $v \neq 0$, $\lambda \in \mathbb{R}_{\max}$, such that

$$Av = \lambda v$$

$$A^k v = \lambda^k v$$

Eigenvector

Find $v \in \mathbb{R}_{\max}^d$, $v \not\equiv -\infty$, $\lambda \in \mathbb{R}_{\max}$, such that

$$\max_{j \in [d]} A_{ij} + v_j = \lambda + v_i$$

$$A^k v = k\lambda + v$$

Theorem (Max-plus spectral theorem, Cuninghame-Green, 61, Gondran & Minoux 77, Cohen et al. 83)

Assume $G(A)$ is strongly connected. Then

- the eigenvalue is unique:

$$\rho_{\max}(A) := \max_{i_1, \dots, i_k} \frac{A_{i_1 i_2} + \dots + A_{i_k i_1}}{k}$$

- Assume WLOG $\rho_{\max}(A) = 0$, then, $\exists \alpha_j \in \mathbb{R} \cup \{-\infty\}$,

$$u = \max_{j \in \text{maximizing circuits}} \alpha_j + A_{\cdot j}^*$$

$A_{ij}^* := \max \text{ weight path arbitrary length } i \rightarrow j.$

- “ $A^{N+c} = \rho_{\max}(A)^c A^N$ ”, $\exists N, c$

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- $A^{N+c} = c\rho_{\max}(A) + A^N, \exists N, c$

The dual linear problem of

$$\min \lambda, A_{ij} + v_j \leq \lambda + u_i \quad \forall i, j$$

is

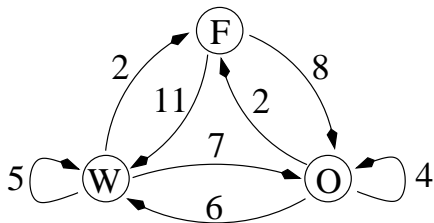
$$\rho(A) = \max_x \sum_{ij} A_{ij} x_{ij}, \quad x_{ij} \geq 0, \quad \sum_j x_{ij} = \sum_j x_{ji}, \quad \sum_{ij} x_{ij} = 1$$

The extreme points of the polytope of circulations are uniform measures supported by elementary circuits.

Complementary slackness shows that v, λ, x optimal iff

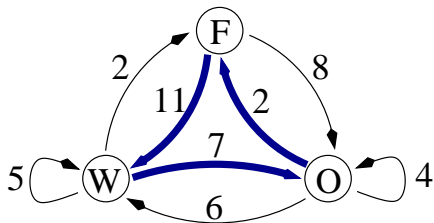
$$x_{ij}(\lambda + u_i - A_{ij} - v_j)$$

Discrete version of maximizing measures.



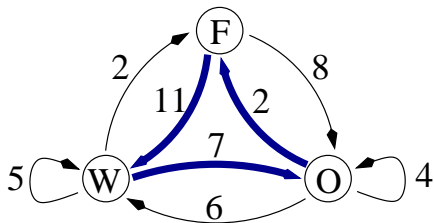
F=fallow (no crop), W=wheat, O=oat, $\rho_{\max}(A) = 20/3$

N. Bacaer, C.R. Acad. d'Agriculture de France, 03



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Actually, **Bacaer** showed that a memory of two years is needed to recover the different historical rotations

The **critical graph** $G^c(A)$ is the union of the maximizing circuits (analogue of Mather and Aubry sets - no difference between them in this discrete case).

Lemma

If i, j are in the same strongly connected component of the critical graph, then $A_{\cdot i}^$ and $A_{\cdot j}^*$ are tropically proportional.*

$$"A^* A^* = A^*"$$

$$\max_k A_{ik}^* + A_{kj}^* = A_{ij}^*$$

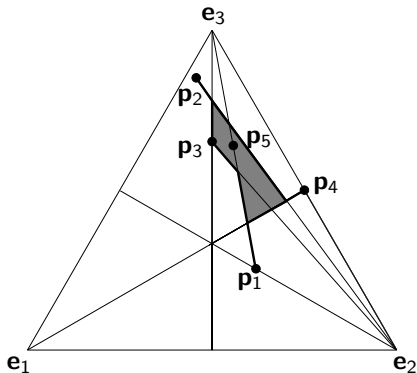
i, j in the same component means $A_{ij}^* + A_{ji}^* = 0$.

$$A_{kj}^* \geq A_{ki}^* + A_{ij}^* \geq A_{kj}^* + A_{ji}^* + A_{ij}^* = A_{kj}^*$$

A vector $u \in C$ is **extreme** if $u = \sup(v, w)$, $v, w \in C$ implies $u = v$ or $u = w$. I.e.,
 $u \in [v, w]$, $v, w \in C \implies u = v$ or $u = w$.

Theorem (Tropical Minkowski-Carathéodory, SG, Katz LAA07; Butkovič, Sergeev, Schneider LAA07; infinite dim Choquet Poncet thesis 11)

Every element of a closed tropical convex set of \mathbb{R}_{\max}^n is the tropical convex combination of at most n extreme points.



Proof.

$$S_i(u) = \{x \in C \mid x \leq u \mid x_i = u_i\}$$

$$\text{Extr } C = \cup_i \text{Min } S_i$$

Proposition

Every $A_{\cdot j}^*$, $j \in G^c(A)$ is *extreme* in the tropical cone $\{v \mid Av = \lambda v\}$.

Cyclicity

WLOG: $\rho(A) = 0$.

The smallest c such that $A^{N+c} = A^N$ for some N (cyclicity) is

$$c = \text{lcm}(\text{cyc}(K_1), \dots, \text{cyc}(K_s))$$

where K_1, \dots, K_s are the strongly connected components of the critical graph, and the cyclicity of a strongly connected component is the **gcd** of the lengths of its circuits.

Cohen, Dubois, Quadrat, Viot 83, Nussbaum 88

Give example at the blackboard.

- If T is a nonexpansive mapping $\mathbb{R}^n \rightarrow \mathbb{R}^n$ with respect to a polyhedral norm, and if T has bounded orbits, then, any orbit converges to a periodic orbit of length bounded by a function of n and of the number of facets of the ball.

Weller, Sine, Nussbaum, Verdyun-Lunel, Scheutzow, Lemmens, ...

- If T is a Shapley operator (order preserving, additively homogeneous) and convex (=1 player), possible orbits lengths are the orders of permutations Akian, SG 03.
- If T is a Shapley operator (2-player), the optimal bound on the length is $\binom{n}{\lfloor n/2 \rfloor}$, the size of a maximal antichain in $\{0, 1\}^n$: Lemmens and Scheutzow, Ergodic Th. and Dyn. Sys.
- If T is sup-norm nonexpansive, Nussbaum conjectured the optimal length to be 2^n .

Spectral projector

WLOG $\rho(A) = 1$, $c = 1$.

$$A^N = A^{N+1} = \dots = P, \quad P = P^2, \quad AP = PA$$

$$P_{ij} = \sup_k A_{ik}^* + A_{kj}^*$$

= **Turnpike theorem** (every long path goes through a maximizing circuit).

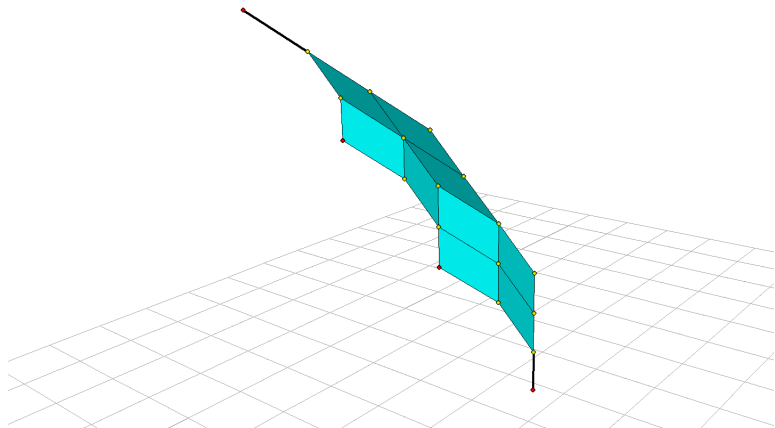
Let K denote the set of critical nodes, $E = \{u \mid Au = u\}$.
The restriction $u \mapsto (u_i)_{i \in K}$ (trace on the projected Aubry set) is an isomorphism, with image

$$\{v \in \mathbb{R}^K \mid v_i - v_j \geq A_{ij}^*, \quad \forall i, j \in K\}$$


= Space of Lipschitz functions for the metric $-A^*$

Note all the tropical convex sets are images of linear projectors. The images of linear projectors arise precisely in this way.

For a Shapley operator (2 player), the tropical convex set $\{u \mid u \leq T(u)\}$ is a polyhedral complex Develin, Sturmfels Doc. Math. 04. Every cell of this complex corresponds to a strategy, and is the image of a linear projector.



The representation of the eigenspace carries over to the infinite dimensional setting.

- Generalizations to kernels appeared in works of **Nussbaum and Mallet-Parret**, under quasi-compactness conditions (essential spectral radius)
-  the existence of a continuous eigenvector is in general a difficult problem.
- Lax-Oleinik semigroups treated in **book by Maslov and Kolokoltsov, Kluwer 97** (typically when the projected Aubry set is finite). Spectral projector written in this context. WKB asymptotics.

Here: abstract boundary theory

Martin boundary, discrete case (Dynkin)

Given P_{xy} Markov kernel, over a discrete infinite set E , find all nonnegative harmonic functions: $u = Pu$.

- 1) Define the Green kernel: $G = P^0 + P + P^2 + \dots$
- 2) The Martin kernel is:

$$K_{xy} = \frac{G_{xy}}{G_{by}}$$

where $b \in E$ is a basepoint.

- 3) Let $\mathcal{K} := \{K_{\cdot y} \mid y \in E\}$
- 4) The Martin space \mathcal{M} is the closure of \mathcal{K} in the product topology.
- 5) The Martin boundary is $\mathcal{B} := \mathcal{M} \setminus \mathcal{K}$.

Theorem (classical Martin representation)

Every harmonic function u can be written as a positive linear combination of functions from the boundary:

$$u = \int_{\mathcal{B}} w \mu(dw) .$$

μ can be chosen to be supported by a subset of \mathcal{B} , the minimal Martin boundary. (We recognise Choquet's theorem!).

Computing the probabilistic Martin boundary is difficult, eg. **Ney and Spitzer 65**, boundary of random walk in \mathbb{Z}^2 is the circle, computing the tropical analogue is much easier!

The max-plus Martin boundary

Akian, SG, Walsh, CDC06, Doc. Math. 09 (Semigroup version),
Ishii, Mitake 07 (PDE version).

Consider the eigenproblem over an arbitrary state space S

$$u_x = \sup_{y \in S} A_{xy} + u_y, \quad \forall x \in S$$

The **Martin kernel** reads: $K_{xy} = A_{xy}^* - A_{by}^*$.

The Martin space \mathcal{M} is the closure of $\mathcal{K} := \{K_{\cdot, y} \mid y \in S\}$ in the product topology (compact, Tychonoff). Martin boundary (set of **horofunctions**) is $\mathcal{B} = \mathcal{M} \setminus \mathcal{K}$.

When $A_{x,y}^* = -d(x,y)$ is the opposite of a metric, recover the construction of the **horoboundary** by Gromov.

The detour metric

$$A^* = "I + A^+", \quad A^+ = "A + A^2 + A^3 + \dots"$$

$$A_{xy}^+ = \sup(A_{xy}, A_{xy}^2, A_{xy}^3, \dots)$$

$$H_{xy}^b = A_{bx}^+ + A_{xy}^+ - A_{by}^+ \quad \text{detour penalty}$$

Extend H^b to the whole Martin space

$$H^b(u, v) = \limsup_{x_d \rightarrow u} \liminf_{y_e \rightarrow v} H_{x_d, y_e}^b$$

where the limsup, inf are taken along nets x_d and y_e converging to u and v in the topology of the Martin space.

The Minimal Martin space is

$$\mathcal{M}^m := \{w \in \mathcal{M} \mid H^b(w, w) = 0\}.$$

Theorem (Max-plus Martin representation Akian, SG, Walsh, CDC06, Doc. Math. 09)

\mathcal{M}^m is the set of extreme elements of $\{u \mid Au = u\}$. Any such u can be written as

$$u = \sup_{w \in \mathcal{M}_m} w + \mu(w), \quad \mu : \mathcal{M}_m \rightarrow \mathbb{R} \cup \{-\infty\} \quad \text{scs}$$

$$\mu_u(w) := \limsup_{x_d \rightarrow w} A_{bx_d}^* + u(x_d)$$

Analogous to max-plus integral representations by Fathi, Siconolfi, Contreras, Ishii, Mitake, in different settings.

If the Martin space is metrisable, then \mathcal{M}_m is precisely the set of **Busemann points** = limits of quasi-geodesics, i.e. of sequences x_1, x_2, \dots such that there exists $\alpha \in \mathbb{R}$

$$A_{bx_k}^* \leq A_{bx_1}^* + A_{x_1x_2} + \dots + A_{x_{k-1}x_k} + \alpha, \quad \forall k$$

Quasi geodesics correspond to almost-sure trajectories of the renormalized H-process of Dynkin.

Lax-Oleinik (continuous time) version in CDC06.

Linear quadratic control - nonquadratic solutions

Hamilton–Jacobi equation

$$\lambda = -|\mathbf{x}|^2 + \frac{1}{4}|\nabla w|^2$$

Maximise reward:

$$- \int_0^T (|\gamma(t)|^2 + |\dot{\gamma}(t)|^2 + \lambda) dt,$$

If $\lambda > 0$, solutions are

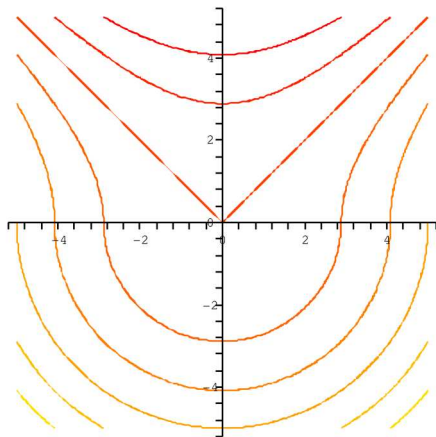
$$w(\mathbf{x}) = \sup_{\mathbf{n}} (\nu(\mathbf{n}) + h_{\mathbf{n}}(\mathbf{x})),$$

where ν is an upper semi–continuous map from the unit vectors to $\mathbb{R} \cup \{-\infty\}$.

When $\lambda = 0$, there is a horofunction for each direction \mathbf{n} :

$$h_{\mathbf{n}}(\mathbf{x}) = \begin{cases} -|\mathbf{x}|^2 + 2(\mathbf{x} \cdot \mathbf{n})^2, & \text{if } \mathbf{x} \cdot \mathbf{n} > 0, \\ -|\mathbf{x}|^2, & \text{otherwise.} \end{cases}$$

The function $-|\mathbf{x}|^2$ is also a horofunction.

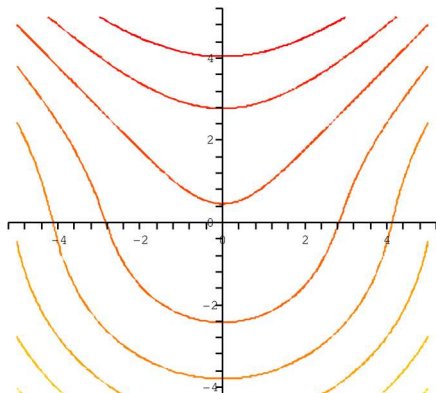


Horospheres of h_n with $n = (0, 1)$.

When $\lambda > 0$: for each direction \mathbf{n} ,

$$h_{\mathbf{n}}(\mathbf{x}) = -\lambda \frac{|\mathbf{x}|^2}{R^2} + \mathbf{x} \cdot \mathbf{n} \frac{\lambda + 2|\mathbf{x}|^2}{R} - \lambda \log \frac{R}{\sqrt{\lambda}},$$

where $R := \sqrt{(\mathbf{x} \cdot \mathbf{n})^2 + \lambda} - \mathbf{x} \cdot \mathbf{n}$.



Back to finite dimension.

The max-plus spectral problem as a limit of the Perron-Frobenius problem

Deformation of the Perron root

Chain of spins (Ising)

$$Z = \sum_{\sigma_1, \dots, \sigma_n \in \Sigma^N} \exp\left(-\sum_{i=1}^N E(\sigma_i, \sigma_{i+1})/T\right), \quad \sigma_{N+1} := \sigma_1$$

$-E(\sigma, \sigma') = H\sigma + J\sigma\sigma'$, $\sigma, \sigma' \in \{\pm 1\}$ (Ising)

$$Z_N = \text{tr } M_T^N, \quad (M_T)_{\sigma\sigma'} = \exp(-E(\sigma, \sigma')/T)$$

$F_N = N^{-1} T \log Z_N \sim T \log \rho(M_T)$ free energy per site,

$T \rightarrow 0$, ground state

$$\epsilon := \exp(-1/T), \quad (M_T)_{\sigma, \sigma'} = \epsilon^{E(\sigma, \sigma')}$$

Similar to perturbation problems, but now, the “Puiseux series” have real exponents (Dirichlet series).

Kingman 61:

$\log \circ \rho \circ \exp$ convex [entrywise exp]

Let $A, B \geq 0$, and $C = A^{(s)} \circ B^{(t)}$, with $s + t = 1, s, t \geq 0$ [entrywise product and exponent] then

$$\rho(C) \leq \rho(A)^s \rho(B)^t .$$

Indeed, $\log \rho(C) = \lim_m \log \|C^m\|/m$ is a pointwise limit of convex functions of $(\log C_{ij})$, for any monotone norm. □

So

$$\rho(A \circ B) \leq \rho(A^{(p)})^{1/p} \rho(B^{(q)})^{1/q} \quad 1/p + 1/q = 1$$

$$\rho(B^{(q)})^{1/q} \rightarrow \max_{i_1, \dots, i_m} (B_{i_1 i_2} \cdots B_{i_{m-1} i_m})^{1/m} =: \rho_\infty(B)$$

Theorem (Friedland 86)

For all $A \in \mathbb{C}^{n \times n}$,

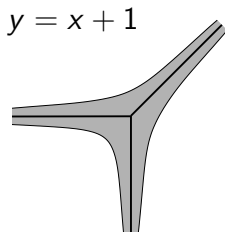
$$\rho(A) \leq \rho(\text{pattern}(A)) \rho_\infty(|A|) \leq n \rho_\infty(|A|)$$

and

$$\rho(A) \geq \rho_\infty(A) \quad \text{if } A_{ij} \geq 0 .$$

Explanation: approximation of an amoeba by its skeleton

$$V \subset (\mathbb{C}^*)^n, A(V) = \{(\log |z_1|, \dots, \log |z_n|) \mid x \in V\}.$$



Cf. Gelfand, Kapranov, Zelevinsky; Passare, Rüllgaard; Purbhoo; Yger.

Limit of the Perron eigenvector. Consider $A^{(p)} = (A_{ij}^p)$, and let $U(p)$ denote the normalized Perron eigenvector of $A^{(p)}$.

Taking $p^{-1} \log$ / passing in the limit in

$$\lambda(p) U_i^p(p) = \sum_j A_{ij}^p U_j^p$$

we get that

$$\lambda + u_i = \max_j \log A_{ij} + u_j$$

where λ and u_j are accumulation points of $p^{-1} \log \lambda(p)$, $\log U_j(p)$, resp.

Which tropical eigenvector is selected?

WLOG $\lambda = \log \rho_\infty(A) = 0$.

Theorem (Akian, Bapat, SG CRAS 1998)

If there is only one SCC of the critical graph with maximal Perron root, then $u_j = (\log A)_{ij}^$, for any j in this class.*

Related work by Lopes, Mohr, Souza, Thieullen.

Give example at the blackboard.

Proof idea. Make diagonal scaling

$$B(p) = \text{diag}(\exp(-pu))A^p \text{diag}(\exp(pu)) .$$

The matrix $B(p)$ has a limit in $[0, 1]^{n \times n}$ as $p \rightarrow \infty$.

We want $B(\infty)$ to have a positive eigenvector. A nonnegative matrix has a positive eigenvector iff the basic classes are exactly the final classes.

For the choice of eigenvector $u = (\log A)_{\cdot j}^*$, this is the case, because the saturation graph

$$\{(k, l) \mid \log A_{kl} + u_l = u_k\}$$

is a river network with sea $\text{SCC}(j)$. Make drawing.

An application: perturbation of eigenvalues

Exercise.

$$\mathcal{A}_\varepsilon = \begin{bmatrix} \varepsilon & 1 & \varepsilon^4 \\ 0 & \varepsilon & \varepsilon^{-2} \\ \varepsilon & \varepsilon^2 & 0 \end{bmatrix},$$

An application: perturbation of eigenvalues

Exercise.

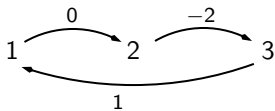
$$\mathcal{A}_\epsilon = \begin{bmatrix} \epsilon & 1 & \epsilon^4 \\ 0 & \epsilon & \epsilon^{-2} \\ \epsilon & \epsilon^2 & 0 \end{bmatrix},$$

Show without computation that the eigenvalues have the following asymptotics as $\epsilon \rightarrow 0$

$$\mathcal{L}_\epsilon^1 \sim \epsilon^{-1/3}, \mathcal{L}_\epsilon^2 \sim j\epsilon^{-1/3}, \mathcal{L}_\epsilon^3 \sim j^2\epsilon^{-1/3}.$$

$$\mathcal{A}_\varepsilon = \begin{bmatrix} \varepsilon & 1 & \varepsilon^4 \\ 0 & \varepsilon & \varepsilon^{-2} \\ \varepsilon & \varepsilon^2 & 0 \end{bmatrix}, \quad A = \begin{bmatrix} 1 & 0 & 4 \\ \infty & 1 & -2 \\ 1 & 2 & \infty \end{bmatrix}.$$

We have $\gamma_1 = -1/3$, corresponding to the critical circuit:



Eigenvalues:

$$\mathcal{L}_\varepsilon^1 \sim \varepsilon^{-1/3}, \mathcal{L}_\varepsilon^2 \sim j\varepsilon^{-1/3}, \mathcal{L}_\varepsilon^3 \sim j^2\varepsilon^{-1/3}.$$

Assume that the entries of \mathcal{A}_ϵ have Puiseux series expansions in ϵ , or even that $\mathcal{A}_\epsilon = a + \epsilon b$, $a, b \in \mathbb{C}^{n \times n}$.

$\mathcal{L}_1, \dots, \mathcal{L}_n$ eigenvalues of \mathcal{A}_ϵ .

$v(s)$: opposite of the smallest exponent of a Puiseux series s .

$\gamma_1 \geq \dots \geq \gamma_n$: tropical eigenvalues of $v(A_\epsilon)$.

Theorem (Akian, Bapat, SG CRAS04, arXiv:0402090)

$$v(\mathcal{L}_1) + \dots + v(\mathcal{L}_n) \leq \gamma_1 + \dots + \gamma_n$$

and equality holds under generic (Lidski-type) conditions.

The maximal tropical eigenvalue γ_1 coincides with the ergodic constant of the one-player game

$$\lambda + u_i = \max_{1 \leq j \leq n} (\text{val}(A_\epsilon)_{ij} + u_j), \forall i$$

λ is the maximal circuit mean.

In general, tropical eigenvalues are non-differentiability points of a parametric optimal assignment problem = Legendre transform of the generic Newton polygon

The (algebraic) **tropical eigenvalues** of a matrix $A \in \mathbb{R}_{\max}^{n \times n}$ are the roots of

$$\text{“per}(A + xI)\text{”}$$

where

$$\text{“per}(M)\text{”} := \sum_{\sigma \in S_n} \prod_{i \in [n]} M_{i\sigma(i)}$$



All geom. eigenvalues λ (“ $Au = \lambda u$ ”) are algebraic eigenvalues, but the converse does not hold. $\rho|_{\max}(A)$ is the max algebraic eigenvalue.

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- Trop. eigs. can be computed in $O(n)$ calls to an optimal assignment solver (Butkovič and Burkard) (not known whether the formal characteristic polynomial can be computed in polynomial time).

Theorem (Kapranov)

If $f(z) = \sum_k f_k z^k \in \mathbb{C}\{\{\epsilon\}\}[z_1, \dots, z_n]$, the closure of the image of $f = 0$ by v is the set of points $x \in \mathbb{R}^n$ at which the maximum

$$\max_k v(f_k) + \langle k, x \rangle$$

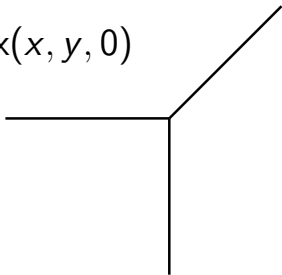
is attained at least twice.

Follows from Puiseux theorem when $n = 1$. Inclusion \subset obvious. Converse: reduction to Puiseux.

When $n = 1$: the set of tropical roots is a zero-dimensional amoeba

Example. $y = x + 1$, $K = \mathbb{C}\{\{\epsilon\}\}$

$\max(x, y, 0)$



Algorithms for games

$$H_i : \max_{1 \leq j \leq n} a_{ij} + x_j \leq \max_{1 \leq k \leq n} b_{ik} + x_k$$

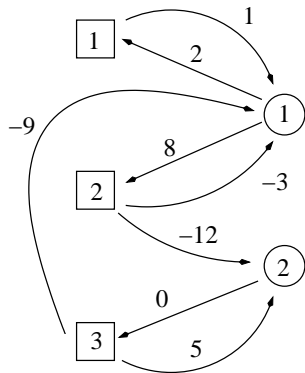
$$[T(x)]_j = \inf_{i \in I} -a_{ij} + \max_{1 \leq k \leq n} b_{ik} + x_k .$$

Interpretation of the game

- State of MIN: variable x_j , $j \in \{1, \dots, n\}$
- State of MAX: half-space H_i , $i \in I$
- In state x_j , Player MIN chooses a tropical half-space H_i with x_j in the LHS
- In state H_i , player MAX chooses a variable x_k at the RHS of H_i
- Payment $-a_{ij} + b_{ik}$.

$$A = \begin{pmatrix} 2 & -\infty \\ 8 & -\infty \\ -\infty & 0 \end{pmatrix}$$

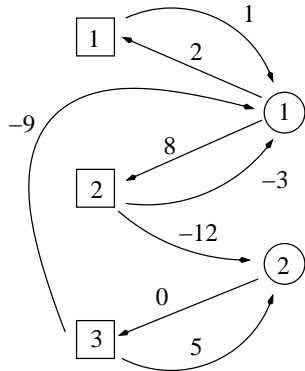
$$B = \begin{pmatrix} 1 & -\infty \\ -3 & -12 \\ -9 & 5 \end{pmatrix}$$



$$2 + x_1 \leq 1 + x_1$$

$$8 + x_1 \leq \max(-3 + x_1, -12 + x_2)$$

$$x_2 \leq \max(-9 + x_1, 5 + x_2)$$



$$\chi(T) = \lim_k v^k / k = (-1, 5)$$

Proposition

If T is nonexpansive and piecewise affine $\mathbb{R}^n \rightarrow \mathbb{R}^n$, the discounted value $v_\alpha = T(\alpha v_\alpha)$ has a Laurent series expansion

$$v_\alpha = \frac{a_{-1}}{1 - \alpha} + a_0 + (1 - \alpha)a_1 + \dots, a_i \in \mathbb{R}^n$$

This is the case for a stochastic game with perfect information and finite action spaces.

Then

$$\chi(T) = \lim_k T^k(0)/k = a_{-1} .$$

- Strategy of MAX $\sigma : \{H_1, \dots, H_m\} \rightarrow \{x_1, \dots, x_n\}$, in state H_i choose coordinate $x_{\sigma(i)}$

Duality theorem (coro of Kohlberg) If finite action spaces, then

$$\chi(T) = \max_{\sigma} \chi(T^{\sigma}) = \min_{\pi} \chi(T_{\pi}) .$$

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- Strategy of MIN $\pi : \{1, \dots, n\} \rightarrow \{1, \dots, m\}$, in state x_j choose hyperplane $H_{\pi(j)}$

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- Strategy of MIN $\pi : \{1, \dots, n\} \rightarrow \{1, \dots, m\}$, in state x_j choose hyperplane $H_{\pi(j)}$
- One player Shapley operators

$$[T^\sigma(x)]_j = \inf_{1 \leq i \leq m} -a_{ij} + b_{i\sigma(i)} + x_{\sigma(i)} .$$

$$[T_\pi(x)]_j = -a_{\pi(j)j} + \max_{1 \leq k \leq n} b_{\pi(j)k} + x_k .$$

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Duality theorem (coro of Kohlberg) If finite action spaces, then

$$\chi(T) = \max_{\sigma} \chi(T^\sigma) = \min_{\pi} \chi(T_\pi) .$$

Every $\chi(T^\sigma)$ and $\chi(T_\pi)$ can be computed in polynomial time.

Proof: Blackwell optimality

For all $x \in \mathbb{R}^n$, we have a selection

$$\exists \sigma, \pi, T(x) = T^\sigma(x) = T_\pi(x) .$$

So for all $0 < \alpha < 1$, the discounted value $v_\alpha = T(\alpha v_\alpha)$ satisfies

$$v_\alpha(T) = \max_{\sigma} v_\alpha(T^\sigma) = \min_{\pi} v_\alpha(T_\pi) .$$

Since χ is the first coefficient of the Laurent series

$$\chi(T) = \max_{\sigma} \chi(T^\sigma) = \min_{\pi} \chi(T_\pi) .$$

σ, π are **Blackwell optimal** if optimal for all $\alpha \in (\bar{\alpha}, 1)$ (exist because the zeros of a Laurent series cant accumulate at 1^-).

Corollary (Condon 92, Zwick and Paterson, TCS 96)

Mean payoff games are in $NP \cap co-NP$.

- I can convince you that $\chi_i(T) \geq 0$ (initial state i is winning) by giving you a strategy σ of MAX such that $\chi_i(T^\sigma) \geq 0$. You can check that in polynomial time by solving a one player game.
- I can convince you that the opposite is true by giving you a strategy π of MIN such that $\chi_i(T_\pi) < 0$. You can also check this in polynomial time.

The class $NP \cap co-NP$ captures the **good characterizations** of **Edmonds**. Evidence that the problem is **not NP-complete**.

- “ $Ax \leq Bx$ ” unfeasible iff $\exists \pi, \bar{\chi}(T_\pi) < 0$.

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- “ $Ax \leq Bx$ ” unfeasible iff $\exists \pi, \bar{\chi}(T_\pi) < 0$.
- “ $Ax \leq Bx$ ” feasible iff $\exists \sigma, \bar{\chi}(T^\sigma) \geq 0$.
- $\exists x \in \mathbb{R}_{\max}^n, Ax \leq Bx?$ is in $\text{NP} \cap \text{co-NP}$

Corollary

Feasibility in tropical linear programming, i.e.,

$$\exists u \in (\mathbb{R} \cup \{-\infty\})^n, \max_j a_{ij} + u_j \leq \max_j b_{ij} + u_j, 1 \leq i \leq p$$

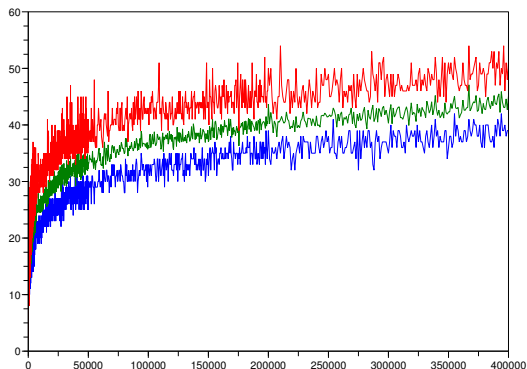
*is polynomial-time equivalent to **mean payoff games**.*

are in $\text{NP} \cap \text{coNP}$: **Zwick, Paterson 96**.

Tropical convex sets are log-limits of classical convex sets: polynomial time solvability of mean payoff games might follow from a **strongly** polynomial-time algorithm in linear programming (**Schewe**).

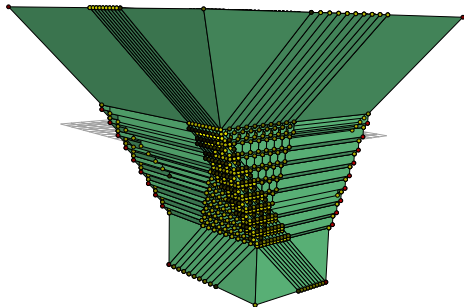
Several pseudo-polynomial algorithms exist for (deterministic) mean payoff games: [Zwick, Paterson TCS96](#). No pseudo-polynomial algorithm seems to be known for stochastic mean payoff game. However, Policy iteration works ([Cochet,SG 06](#)), - based on a tropical idea = spectral projectors - ; alternative algorithm by [Boros, Gurvich, Elbassioni, Makino, . . .](#)

Policy iteration for games scales well in practice. $\#$
iterations / $\#$ nodes



However, **Friedmann LICS 10** showed that policy iteration for games can be exponential.

Intersection of 10 affine tropical hyperplanes in dimension 3, only 24 vertices, but 1215 pseudo-vertices.



Tropical double description [Allamigeon, SG, Goubault](#).
Efficient implementation in TPLib/caml by [Allamigeon](#).

Concluding remarks

- Tropical algebra \sim discrete version of Weak KAM
- Tropical convex cones arises when considering spaces of weak KAM solutions (1-player), or sub/super solutions.
- Combinatorial properties in the discrete case (lengths of periodic orbits)
- Thinking tropical brings “complex” perspective on Lax-Oleinik semigroups (not just one eigenvalue)
- Relation between ergodic problem and optimal assignment appears in the discrete case (the eigenvalues are nondifferentiability points of an optimal assignment problem), is there a PDE analogue (relation with mass transport problem)?
- Tropical algebra is fun!

Thank you