

Elliptic curve V

Divisors and Miller's algorithm

I) Motivations:

Goal: compute pairings

The Weil pairing $e_{w,r}: E[l] \times E[l] \rightarrow \mu_l$
 $P, Q \longrightarrow \frac{f_{l,P}(Q)}{f_{l,Q}(P)}$

$f_{l,P}$ is the unique rational function in $k(E)$ with a zero of order l at P and a pole of order l at O_E , which is normalized at infinity

Goal: construct $f_{l,P}$. More generally, given poles and zeroes with multiplicities, construct a function with these multiplicities if it exists.

II) Divisors on curves:

Let \mathcal{C}/k be a smooth projective curve. Recall that for each $P \in \mathcal{C}(k)$, there is a corresponding valuation v_P on $k(\mathcal{C})$.

If $f \in k(\mathcal{C})$ is a rational function on \mathcal{C} .

Define $\text{div}(f) = \sum_{P \in \mathcal{C}(k)} v_P(f) P$ a formal sum of points.

Example: $f = \frac{x^2(x-2)}{x-1}$ in $k(\mathbb{P}^1) = k(x)$

$\text{div}(f) = 2(x=0) + 1(x=2) - 1(x=1) - 2(x=\infty)$

Lemma: if $f \in k(\mathcal{C})$, it only has a finite number of poles and zeroes, so $\text{div}(f)$ only involve a finite normal sum.

Definition: A divisor D on \mathcal{C} is a finite formal sum of points:

$$D = \sum_{P \in \mathcal{C}(k)} m_P(P), \text{ where } m_P = 0 \text{ for all but a finite number of points.}$$

Definition: $\text{deg}(D) = \sum m_P$

Theorem: 1) If $f \in k(\mathcal{C})$, $\text{deg } \text{div}(f) = 0$

2) If $\text{div}(f) = 0$ i.e. if f has no pole or zeroes, then f is constant: $f = \alpha, \alpha \in k^*$

Corollary: if $\text{div}(f) = \text{div}(g)$ then $\text{div}(f/g) = \text{div}(f) - \text{div}(g) = 0$
then $f = \alpha g, \alpha \in k^*$

Definition: * D is linearly equivalent to 0 if $D = \text{div}(f)$

* D and D' are linearly equivalent if $D - D'$ is linearly equivalent to 0 , i.e. $D = D' + \text{div}(f)$

* $\text{Pic}^0(\mathcal{C}) = \{ \text{Divisors of } \mathcal{C} \text{ of degree } 0 \}$ modulo linear equivalence
(this is a group since we can sum divisors and this is compatible with equivalence)

III) Divisors on \mathbb{P}^1

Let $f \in k(\mathbb{P}^1) = k(x)$. Over \bar{k} , $f = \frac{\prod (x - \alpha_i)^{m_i}}{\prod (x - \beta_i)^{n_i}} \cdot x^r$

$$\text{deg } f = \sum m_i - \sum n_i$$

$$\text{div}(f) = \sum m_i (\alpha_i) - \sum n_i (\beta_i) - \text{deg } f (\infty)$$

$$\text{deg } \text{div}(f) = 0$$

* If f has no poles and zeroes, over k its denominator and numerator have degree 0, so $f = \frac{a}{b} = \gamma \in k^*$

* Let $D = \sum n_i (\alpha_i)$ with $\deg(D) = 0$, so $\sum n_i = 0$. Does D come from a rational function?

$$D = \sum_{\alpha_i \neq \infty} n_i (\alpha_i) - (\sum n_i) (\infty)$$

$$f_D = \prod (x - \alpha_i)^{n_i}$$

$$\text{div}(f_D) = D$$

There is no obstruction i.e. $\text{Pic}^0(\mathbb{P}^1) = 0$

* Normalisation: we say that f_D is normalised (at infinity) if

$$-\text{div}(f_D) = D$$

$$-f_D = \frac{\sum_{i=0}^n a_i x^i}{\sum_{i=0}^m b_i x^i} \quad \text{then} \quad \frac{a_n}{b_m} = 1$$

The quotient of the coefficients of the highest degree is equal to 1

Remember that $\Pi_\infty = \frac{1}{x}$ is a uniformiser at ∞

If $f \in k(x)$, $\frac{f}{\Pi_\infty^{\deg(f)}}$ is well defined at ∞

$$f \text{ is normalised at } \infty \iff \left(\frac{f}{\Pi_\infty^{\deg(f)}} \right) (\infty) = 1$$

IV) Divisors on E:

Lemma: If $f \in k(E)$, $\deg \operatorname{div} f = 0$

Proof: If $\operatorname{div} f = \sum n_P(P)$
 $\operatorname{div} \bar{f} = \sum n_{\bar{P}}(\bar{P})$

$$2 \deg \operatorname{div} f = \deg \operatorname{div} f + \deg \operatorname{div} \bar{f} = \deg \operatorname{div} (f\bar{f}) = \deg \operatorname{div} (Nf)$$

$Nf \in k(X)$

$$\operatorname{div} Nf = \sum_{\text{Not Weierstrass}} v_P(Nf)(P) + \sum_{\text{Weierstrass}} v_P(Nf)(P)$$

* If P not Weierstrass, $v_P(Nf) = v_{x=\alpha_P}(Nf)$ because $x - \alpha_P$ is a uniformizer.
But $-\bar{P}$ has the same x .

$$v_P(Nf) + v_{-\bar{P}}(Nf) = 2 v_{x=\alpha_P}(Nf)$$

* If $P = (\alpha_P, 0)$ is a Weierstrass, it is the only point with x -coordinate α_P , but y is a uniformizer, and $v_P(x - \alpha_P) = 2$

$$\text{So } v_P(Nf) = 2 v_{x=\alpha_P}(Nf)$$

* Conclusion: $\deg \operatorname{div}_E Nf = 2 \deg \operatorname{div}_{\mathbb{P}^1} Nf = 0$

Theorem: Let $D = \sum n_P(P) \in \operatorname{Div}^0(E)$. Then $D \sim 0$ if and only if $\sum n_P P = 0$ in $E(\mathbb{K})$

Corollary: Let $\Sigma_D = \sum n_P P \in E(\mathbb{K})$, $D \sim (\Sigma_D) - (O_E)$

Proof: $D - (\Sigma_D) + (O_E)$ has degree 0, and $\sum n_P P - \Sigma_D = O_E$ by definition

Reformulation: $\operatorname{Pic}^0(E) \cong E(\mathbb{K})$: "algebraic interpretation of the group law"

I) Miller's algorithm:

Goal: Given D such that $\deg D = 0$ and $\sum D_i = O_E$
 construct $f_D \llcorner$ normalized at infinity \gg such that $\text{div } f_D = D$

* $\Pi_{O_E} = \frac{x}{y}$ is a uniformizer at O_E
 f_D is normalized at infinity if $\left(\frac{f_D}{\prod_{O_E} v_{O_E}(f_D)} \right) (O_E) = 1$

if and only if α, β are the coeffs of highest degree of the numerator and denominator, $\frac{\alpha}{\beta} = 1$

* Fundamental case: $D = (P) + (Q) - (P+Q) - (O_E)$

Definition: $\mu_{P,Q}$ is the normalized function with divisor $(P) + (Q) - (P+Q) - (O_E)$

* Reduction: $D = (P_1) + (P_2) + (P_3) + \dots$

$D = \text{div } \mu_{P_1, P_2} + (P_1 + P_2) + (O_E) + (P_3) + \dots = \text{div } \mu_{P_1, P_2} + \text{div } \mu_{P_1 + P_2, P_3} + (P_1 + P_2 + P_3) + \dots$

Note: $\text{div } \mu_{P, -P} = (P) + (-P) - 2(O_E)$

So $-(P) + (O_E) \sim (-P) - (O_E)$

Example: $D = 5(P) - 5(O_E)$ where $P \in E[5]$

$D = \text{div } \mu_{P, P} + (2P) + (3P) - 4(O_E) = \text{div } \mu_{P, P} + \text{div } \mu_{P, 2P} + (3P) + 2(P) - 3(O_E)$
 $= \text{div } \mu_{P, P} + \text{div } \mu_{P, 2P} + \text{div } \mu_{P, 3P} + (4P) + (P) - 2(O_E)$

$= \text{div } \mu_{P, P} + \text{div } \mu_{P, 2P} + \text{div } \mu_{P, 3P} + \text{div } \mu_{P, 4P} + (5P) - (O_E)$
 as $P \in E[5]$

$\Rightarrow f_D = \mu_{P, P} \mu_{P, 2P} \mu_{P, 3P} \mu_{P, 4P}$ is normalized because the $\mu_{P, \dots}$ are

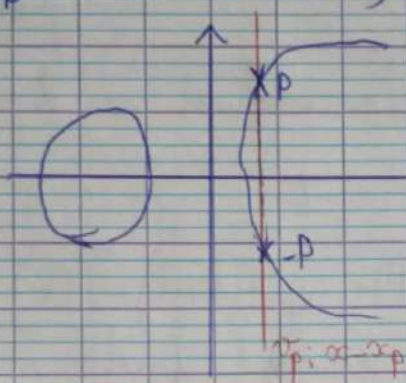
* More generally, using the $\mu_{P, Q}$ we can always reduce D to $n \cdot (P) - (O_E)$

* Conversely, if $D = (P) - (O_E)$ was ~ 0 with $P \neq O_E$, then the function f_D would induce an isomorphism $f_D: E \cong \mathbb{P}^1$ but this is impossible since they don't have the same genus.

* Constructing $\mu_{P,Q}$

a) $Q = -P / \mu_{P,-P}$

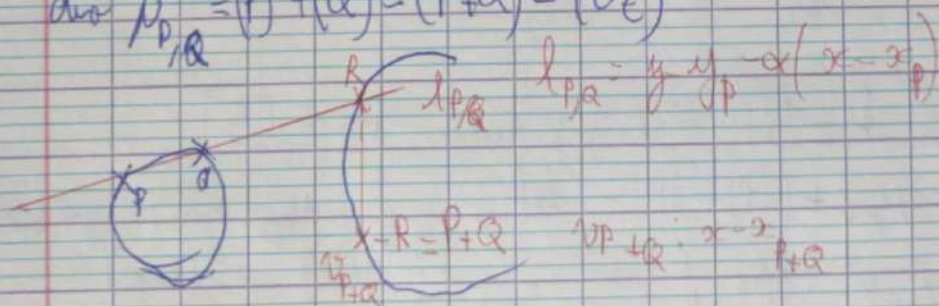
$$\text{div } \mu_{P,-P} = (P) + (-P) - 2(O_E)$$



Lemma: $\mu_{P,-P} = x - x_p$

b) $Q \neq -P$

$$\text{div } \mu_{P,Q} = (P) + (Q) - (P+Q) - (O_E)$$



$$\text{div } l_{P,Q} = (P) + (Q) + (R) - 3(O_E)$$

$$\Rightarrow f_D = \mu_{P,-P} \mu_{P,2P} \mu_{P,3P} \mu_{P,4P}$$

$$\text{div } \nu_{P+Q} = (P+Q) + (R) - 2(\infty)$$

$$M_{P,Q} = \frac{l_{P,Q}}{v_{P+Q}} = \frac{y - y_P - \alpha(x - x_P)}{\alpha - x_{P+Q}}$$

it is normalized

III) Miller's algorithm:

Let $P \in E[l]$, $D = l(P) - l(O_E) \sim 0$

Definition: $f_{l,P}$ the normalized function with this division

Reduction: $(P) + (P) \sim (2P) + (O_E)$ via $\mu_{P,P}$
 $(2P) + (P) \sim (3P) + (O_E)$ via $\mu_{2P,P}$
 $(3P) + (P) \sim (4P) + (O_E)$ via $\mu_{3P,P}$

$\Delta \Rightarrow O(l)$ Too slow

Solution: use the double and add method.

Informally: $(P) + (P) \sim (2P) + (O_E)$
 $(2P) + (2P) \sim (4P) + (O_E)$
 $(4P) + (4P) \sim (8P) + (O_E)$

Rigorously: For any P , define $f_{l,P}$ such that it is normalized and
 $\text{div } f_{l,P} = l(P) - l(P) - (l-1)(O_E)$

Exercise:

- $f_{l,P}$ exists
- If $P \in E[l]$ this is the same function as before
- $f_{l_1+l_2, P} = f_{l_1, P} f_{l_2, P} M_{l_1, P, l_2, P}$

Hint: check that the divisions are the same \Rightarrow the functions are equal since they are normalized.

Coxsley: Double and add

$$\text{Double: } f_{2l,p} = f_{l,p}^2 \mu_{2l,p}$$

$$\text{Add: } f_{l+1,p} = f_{l,p} \mu_{l,p}$$

$$\text{Base: } f_{0,p} = f_{1,p} = 1$$

Miller's algo: Input: $l, p \in \mathbb{E}[l], Q \in \mathbb{E}$
Output: $f_{l,p}(Q)$

- Write $l = \sum_{i=0}^n b_i 2^i$ $T=0$ $f=1$

- For each bit b_i from left to right

$$\text{- Double: } f = f^2 \mu_{T,p}(Q) \\ T = 2T$$

$$\text{- If } b_i = 1 \text{ Add: } f = f \times \mu_{T,p}(Q) \\ T = T + p$$

Return f .