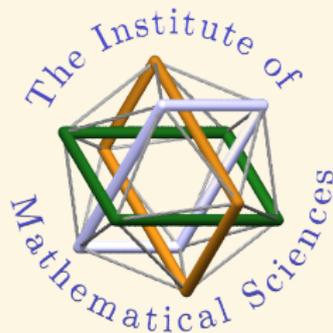


# Optimal pairings on abelian varieties

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# Outline

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## The Weil pairing on elliptic curves

- Let  $E: y^2 = x^3 + ax + b$  be an elliptic curve over a field  $k$  ( $\text{char } k \neq 2, 3$ ,  $4a^3 + 27b^2 \neq 0$ .)
- Let  $P, Q \in E[\ell]$  be points of  $\ell$ -torsion.
- Let  $f_P$  be a function associated to the principal divisor  $\ell(P) - \ell(0)$ , and  $f_Q$  to  $\ell(Q) - \ell(0)$ . We define:

$$e_{W,\ell}(P, Q) = \frac{f_P((Q) - (0))}{f_Q((P) - (0))}.$$

- The application  $e_{W,\ell} : E[\ell] \times E[\ell] \rightarrow \mu_\ell(\bar{k})$  is a non degenerate pairing: the Weil pairing.

### Definition (Embedding degree)

If  $E$  is defined over a finite field  $\mathbb{F}_q$ , the Weil pairing has image in  $\mu_\ell(\bar{\mathbb{F}}_q) \subset \mathbb{F}_{q^d}^*$  where  $d$  is the **embedding degree**, the smallest number such that  $\ell \mid q^d - 1$ .

# The Tate pairing on elliptic curves over $\mathbb{F}_q$

## Definition

The Tate pairing is a non degenerate bilinear application given by

$$e_T: E_0[\ell] \times E(\mathbb{F}_q)/\ell E(\mathbb{F}_q) \longrightarrow \mathbb{F}_{q^d}^*/\mathbb{F}_{q^d}^{*\ell} .$$

$$(P, Q) \longmapsto f_P((Q) - (0))$$

where

$$E_0[\ell] = \{P \in E[\ell](\mathbb{F}_{q^d}) \mid \pi(P) = [q]P\}.$$

- On  $\mathbb{F}_{q^d}$ , the Tate pairing is a non degenerate pairing

$$e_T: E[\ell](\mathbb{F}_{q^d}) \times E(\mathbb{F}_{q^d})/\ell E(\mathbb{F}_{q^d}) \rightarrow \mathbb{F}_{q^d}^*/\mathbb{F}_{q^d}^{*\ell} \simeq \mu_\ell;$$

- If  $\ell^2 \nmid E(\mathbb{F}_{q^d})$  then  $E(\mathbb{F}_{q^d})/\ell E(\mathbb{F}_{q^d}) \simeq E[\ell](\mathbb{F}_{q^d})$ ;
- We normalise the Tate pairing by going to the power of  $(q^d - 1)/\ell$ .

# Miller's functions

- We need to compute the functions  $f_{\ell,P}$  and  $f_{\ell,Q}$ . More generally, we define the Miller's functions:

## Definition

Let  $\lambda \in \mathbb{N}$  and  $X \in E[\ell]$ , we define  $f_{\lambda,X} \in k(E)$  to be a function thus that:

$$(f_{\lambda,X}) = \lambda(X) - ([\lambda]X) - (\lambda - 1)(0).$$

- We want to compute (for instance)  $f_{\ell,P}((Q) - (0))$ .

# Miller's algorithm

- The key idea in Miller's algorithm is that

$$f_{\lambda+\mu, X} = f_{\lambda, X} f_{\mu, X} f_{\lambda, \mu, X}$$

where  $f_{\lambda, \mu, X}$  is a function associated to the divisor

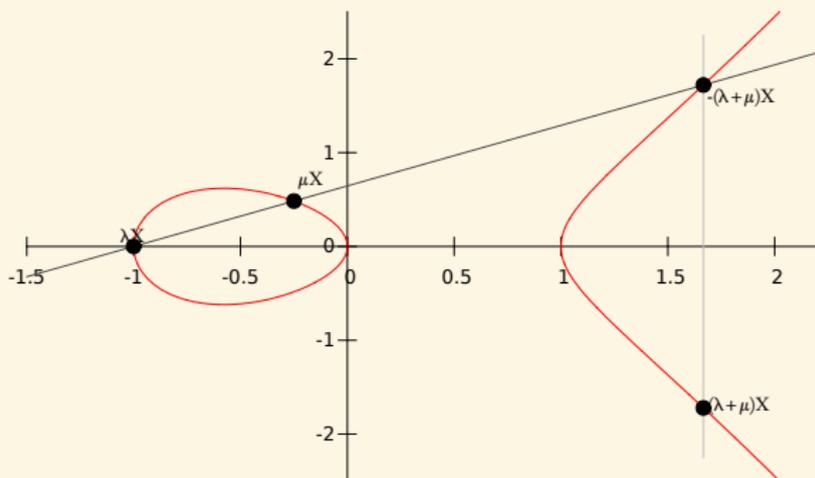
$$([\lambda]X) + ([\mu]X) - ([\lambda + \mu]X) - (0).$$

- We can compute  $f_{\lambda, \mu, X}$  using the addition law in  $E$ : if  $[\lambda]X = (x_1, y_1)$  and  $[\mu]X = (x_2, y_2)$  and  $\alpha = (y_1 - y_2)/(x_1 - x_2)$ , we have

$$f_{\lambda, \mu, X} = \frac{y - \alpha(x - x_1) - y_1}{x + (x_1 + x_2) - \alpha^2}.$$

# Miller's algorithm for elliptic curves

$$[\lambda]X = (x_1, y_1) \quad [\mu]X = (x_2, y_2)$$



$$f_{\lambda, \mu, X} = \frac{y - \alpha(x - x_1) - y_1}{x + (x_1 + x_2) - \alpha^2}.$$

# Miller's algorithm for the Tate pairing on elliptic curves

## Algorithm (Computing the Tate pairing)

Input:  $\ell \in \mathbb{N}$ ,  $P = (x_1, y_1) \in E[\ell](\mathbb{F}_q)$ ,  $Q = (x_2, y_2) \in E(\mathbb{F}_{q^d})$ .

Output:  $e_T(P, Q)$ .

- ① Compute the binary decomposition:  $\ell := \sum_{i=0}^l b_i 2^i$ . Let  $T = P, f_1 = 1, f_2 = 1$ .
- ② For  $i$  in  $[l..0]$  compute
  - ①  $\alpha$ , the slope of the tangent of  $E$  at  $T$ .
  - ②  $T = 2T$ .  $T = (x_3, y_3)$ .
  - ③  $f_1 = f_1^2 (y_2 - \alpha(x_2 - x_3) - y_3)$ ,  $f_2 = f_2^2 (x_2 + (x_1 + x_3) - \alpha^2)$ .
  - ④ If  $b_i = 1$ , then compute
    - ①  $\alpha$ , the slope of the line going through  $P$  and  $T$ .
    - ②  $T = T + Q$ .  $T = (x_3, y_3)$ .
    - ③  $f_1 = f_1^2 (y_2 - \alpha(x_2 - x_3) - y_3)$ ,  $f_2 = f_2 (x_2 + (x_1 + x_3) - \alpha^2)$ .

Return

$$\left( \frac{f_1}{f_2} \right)^{\frac{q^d - 1}{\ell}}.$$

## Miller's algorithm on Jacobians

- Let  $P \in \text{Jac}(C)[\ell]$  and  $D_P$  a divisor on  $C$  representing  $P$ ;
- By definition of  $\text{Jac}(C)$ ,  $\ell D_P$  corresponds to a principal divisor  $(f_{\ell,P})$  on  $C$ ;
- The same formulas as for elliptic curve define the Weil and Tate-Lichtenbaum pairings:

$$e_W(P, Q) = f_{\ell,P}(D_Q) / f_{\ell,Q}(D_P)$$

$$e_T(P, Q) = f_{\ell,P}(D_Q).$$

- A key ingredient for evaluating  $f_P(D_Q)$  comes from Weil's reciprocity theorem.

### Theorem (Weil)

Let  $D_1$  and  $D_2$  be two divisors with disjoint support linearly equivalent to  $(0)$  on a smooth curve  $C$ . Then

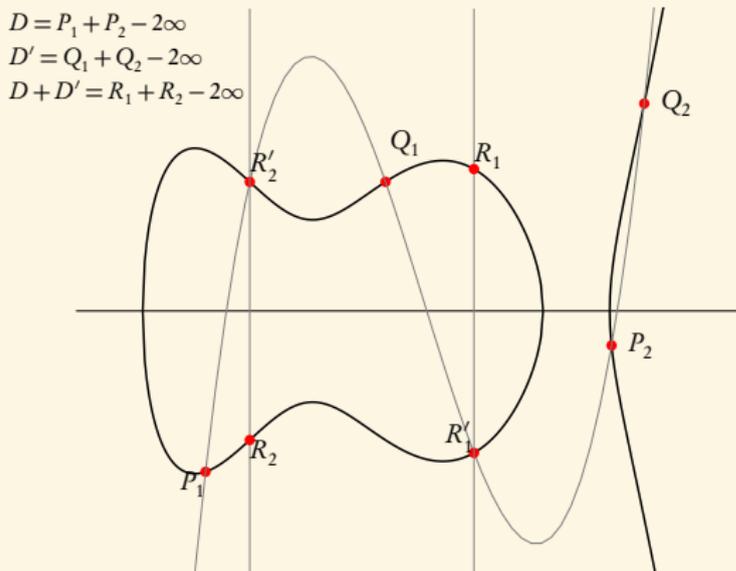
$$f_{D_1}(D_2) = f_{D_2}(D_1).$$

## Miller's algorithm on Jacobians of genus 2 curves

- The extension of Miller's algorithm to Jacobians is “straightforward”;
- For instance if  $g = 2$ , the function  $f_{\lambda, \mu, P}$  is of the form

$$\frac{y - l(x)}{(x - x_1)(x - x_2)}$$

where  $l$  is of degree 3.



# Abelian varieties

## Definition

An **Abelian variety** is a complete connected group variety over a base field  $k$ .

- Abelian variety = **points** on a projective space (locus of homogeneous polynomials) + an abelian group law given by **rational functions**.

## Example

- Elliptic curves = Abelian varieties of dimension 1;
- If  $C$  is a (smooth) curve of genus  $g$ , its Jacobian is an abelian variety of dimension  $g$ ;
- In dimension  $g \geq 4$ , not every abelian variety is a Jacobian.

## The Weil-Cartier pairing

- Let  $f: A \rightarrow B$  be a separable isogeny with kernel  $K$  between two abelian varieties defined over  $k$ ;
- The isogeny  $f$  and its dual  $\hat{f}$  fit into the diagram

$$\begin{array}{ccccccccc}
 0 & \longrightarrow & K & \longrightarrow & A & \xrightarrow{f} & B & \longrightarrow & 0 \\
 & & & & & & & & \\
 & & & & \hat{A} & \xleftarrow{\hat{f}} & \hat{B} & \longleftarrow & \hat{K} & \longleftarrow & 0
 \end{array}$$

- Since  $\hat{K}$  is the Cartier dual of  $K$  we have a non degenerate pairing  $e_f: K \times \hat{K} \rightarrow \mathbb{G}_m$ ;
- Unravelling the identification, we can compute the Weil-Cartier pairing as follows:
  - If  $Q \in \hat{K}(\bar{k})$ ,  $Q$  defines a divisor  $D_Q$  on  $B$ ;
  - $\hat{f}(Q) = 0$  means that  $f^*D_Q$  is equal to a principal divisor ( $g_Q$ ) on  $A$ ;
  - $e_f(P, Q) = g_Q(x)/g_Q(x+P)$ . (This last function being constant in its definition domain).
- The Weil pairing  $e_{W,\ell}$  is the pairing associated to the isogeny  $[\ell]: A \rightarrow A$

$$e_{W,\ell}: A[\ell] \times \hat{A}[\ell] \rightarrow \mu_\ell.$$

## Reformulation

$$\begin{array}{ccc}
 f^*D_Q & \xrightarrow{\psi_Q} & \mathcal{O}_A \\
 \downarrow \psi_P & & \parallel e_f(P, Q) \\
 \tau_P^* f^* D_Q & \xrightarrow{\tau_P^* \psi_Q} & \tau_P^* \mathcal{O}_A
 \end{array}$$

( $\psi_P$  is normalized via  $A(P) \simeq A(0)$ .)

- Since  $f^*D_Q$  is trivial, by descent theory  $D_Q$  is the quotient of  $A \times \mathbb{A}^1$  by an action of  $K$ :

$$g_x \cdot (t, \lambda) = (t + x, \chi_Q(x)\lambda)$$

where  $\chi_Q$  is a character on  $K$ ;

$$e_f(P, Q) = \chi_Q(P).$$

# Polarizations

If  $\mathcal{L}$  is an ample line bundle, the polarization  $\varphi_{\mathcal{L}}$  is a morphism  $A \rightarrow \hat{A}, x \mapsto t_x^* \mathcal{L} \otimes \mathcal{L}^{-1}$ .

## Definition (Weil pairing)

Let  $\mathcal{L}$  be a principal polarization on  $A$ . The (polarized) Weil pairing  $e_{W,\mathcal{L},\ell}$  is the pairing

$$\begin{aligned} e_{W,\mathcal{L},\ell} : A[\ell] \times A[\ell] &\longrightarrow \mu_{\ell} \\ (P, Q) &\longmapsto e_{W,\ell}(P, \varphi_{\mathcal{L}}(Q)) \end{aligned} .$$

associated to the polarization  $\varphi_{\mathcal{L}^{\ell}}$ :

$$A \xrightarrow{[\ell]} A \xrightarrow{\mathcal{L}} \hat{A}$$

## The commutator pairing

- In general for an ample line bundle  $\mathcal{L}$ , the polarization  $\varphi_{\mathcal{L}}$  gives an isogeny

$$0 \longrightarrow K(\mathcal{L}) \longrightarrow A \longrightarrow \hat{A} \longrightarrow 0$$

and thus a pairing

$$e_{\varphi} : K(\mathcal{L}) \times K(\mathcal{L}) \rightarrow \mathbb{G}_m.$$

- The following diagram is commutative up to a multiplication by  $e_{\varphi}(P, Q)$ :

$$\begin{array}{ccc}
 \mathcal{L} & \xrightarrow{\psi_P} & \tau_P^* \mathcal{L} \\
 \downarrow \psi_Q & & \downarrow \tau_P^* \psi_Q \\
 \tau_Q^* \mathcal{L} & \xrightarrow{\tau_Q^* \psi_P} & \tau_{P+Q}^* \mathcal{L}
 \end{array}$$

## The commutator pairing

- The Theta group  $G(\mathcal{L})$  is the group  $\{(x, \psi_x)\}$  where  $x \in K(\mathcal{L})$  and  $\psi_x$  is an isomorphism

$$\psi_x : \mathcal{L} \rightarrow \tau_x^* \mathcal{L}$$

The composition is given by  $(y, \psi_y) \cdot (x, \psi_x) = (y + x, \tau_x^* \psi_y \circ \psi_x)$ .

- $G(\mathcal{L})$  is an Heisenberg group:

$$0 \longrightarrow k^* \longrightarrow G(\mathcal{L}) \longrightarrow K(\mathcal{L}) \longrightarrow 0$$

- Let  $g_P = (P, \psi_P) \in G(\mathcal{L})$  and  $g_Q = (Q, \psi_Q) \in G(\mathcal{L})$ ,

$$e_{\mathcal{L}}(P, Q) = g_P g_Q g_P^{-1} g_Q^{-1};$$

- If  $\psi : K(\mathcal{L}) \times K(\mathcal{L}) \rightarrow k^*$  is the 2-cocycle associated to  $G(\mathcal{L})$ , we also have

$$e_{\mathcal{L}}(P, Q) = \frac{\psi(P, Q)}{\psi(Q, P)}.$$

## Kummer exact sequence

- The exact sequence

$$1 \rightarrow \mu_\ell \rightarrow \bar{k}^* \rightarrow \bar{k}^* \rightarrow 1$$

induces a connecting map

$$\delta : \bar{k}^* / \bar{k}^{*\ell} \simeq H^1(k, \mu_\ell)$$

(the isomorphism comes from Hilbert 90:  $H^1(k, k^*) = 0$ ).

- Thus for a finite field  $k = \mathbb{F}_q$

$$\mathbb{F}_{q^d}^* / \mathbb{F}_{q^d}^{*\ell} \simeq H^1(\mathbb{F}_{q^d}, \mu_\ell) \simeq \mu_\ell(\mathbb{F}_{q^d});$$

- The isomorphism is given by the exponentiation  $x \mapsto x^{\frac{q^d-1}{\ell}}$ .

## The Tate-Cartier pairing on abelian varieties over finite fields

- Let  $f: A \rightarrow B$  be an isogeny with  $\text{Ker } f \subset A[\ell]$ ;
- From the exact sequence

$$0 \rightarrow \text{Ker } f \rightarrow A \rightarrow B \rightarrow 0$$

we get from Galois cohomology a connecting morphism

$$\delta: A(\mathbb{F}_{q^d})/f(B(\mathbb{F}_{q^d})) \simeq H^1(\mathbb{F}_{q^d}, \text{Ker } f)$$

(this is an isomorphism since  $H^1(\mathbb{F}_{q^d}, A) = 0$  for an abelian variety over a finite field);

- Composing with the Weil-Cartier pairing, we get a bilinear application

$$\text{Ker } \hat{f}(\mathbb{F}_{q^d}) \times A(\mathbb{F}_{q^d})/f(B(\mathbb{F}_{q^d})) \rightarrow H^1(\mathbb{F}_{q^d}, \mu_\ell) \simeq \mathbb{F}_{q^d}^*/\mathbb{F}_{q^d}^{*\ell} \simeq \mu_\ell;$$

- Explicitly, if  $P \in \text{Ker } \hat{f}(\mathbb{F}_{q^d})$  and  $Q \in A(\mathbb{F}_{q^d})$  then the (reduced) Tate pairing is given by

$$e_T(P, Q) = e_W(\pi^d(Q_0) - Q_0, P)$$

where  $Q_0 \in A$  is any point such that  $Q = f(Q_0)$  and  $\pi$  is the Frobenius of  $\mathbb{F}_q$ ;

# The Tate-Cartier pairing on abelian varieties over finite fields

## Theorem

### The Tate pairing

$$\text{Ker } \hat{f}(\mathbb{F}_{q^d}) \times A(\mathbb{F}_{q^d})/f(B(\mathbb{F}_{q^d})) \rightarrow H^1(\mathbb{F}_{q^d}, \mu_\ell) \simeq \mathbb{F}_{q^d}^*/\mathbb{F}_{q^d}^{*\ell} \simeq \mu_\ell$$

is non degenerate.

## Proof.

We have canonically

$$\begin{aligned} \text{Ker } \hat{f}(\mathbb{F}_{q^d}) &= \text{Hom}(\text{Ker } f, \mathbb{G}_m)^{\text{Gal}(\overline{\mathbb{F}}_{q^d}/\mathbb{F}_{q^d})} \\ &= \text{Hom}(\text{Ker } f/(\pi^d - 1), \mathbb{F}_{q^d}^*) \\ &= \text{Hom}(H^1(\mathbb{F}_{q^d}, \text{Ker } f), \mathbb{F}_{q^d}^*) \end{aligned}$$

and

$$A(\mathbb{F}_{q^d})/f(B(\mathbb{F}_{q^d})) \simeq H^1(\mathbb{F}_{q^d}, \text{Ker } f).$$



## Cycles and Lang reciprocity

- Let  $(A, \Theta)$  be a principally polarized abelian variety;
- To a degree 0 cycle  $\sum n_i(P_i)$  on  $A$ , we can associate the divisor  $\sum n_i t_{P_i}^* \Theta$  on  $A$ ;
- The cycle  $\sum n_i(P_i)$  corresponds to a trivial divisor iff  $\sum n_i P_i = 0$  in  $A$ ;
- If  $f$  is a function on  $A$  and  $D = \sum(P_i)$  a cycle whose support does not contain a zero or pole of  $f$ , we let

$$f(D) = \prod f(P_i)^{n_i}.$$

(In the following, when we write  $f(D)$  we will always assume that we are in this situation.)

### Theorem (Lang [Lan58])

Let  $D_1$  and  $D_2$  be two cycles equivalent to 0, and  $f_{D_1}$  and  $f_{D_2}$  be the corresponding functions on  $A$ . Then

$$f_{D_1}(D_2) = f_{D_2}(D_1)$$

# The Weil and Tate pairings on abelian varieties

## Theorem

Let  $P, Q \in A[\ell]$ . Let  $D_P$  and  $D_Q$  be two cycles equivalent to  $(P) - (0)$  and  $(Q) - (0)$ . The Weil pairing is given by

$$e_W(P, Q) = \frac{f_{\ell D_P}(D_Q)}{f_{\ell D_Q}(D_P)}.$$

## Theorem

Let  $P \in A[\ell](\mathbb{F}_{q^d})$  and  $Q \in A(\mathbb{F}_{q^d})$ , and let  $D_P$  and  $D_Q$  be two cycles equivalent to  $(P) - (0)$  and  $(Q) - (0)$ . The (non reduced) Tate pairing is given by

$$e_T(P, Q) = f_{\ell D_P}(D_Q).$$

## Cryptographic usage of pairings on abelian varieties

- The Weil pairing was first used to transfer the DLP from an elliptic curve to  $\mathbb{F}_{q^d}^*$  (the MOV attack [MOV91]);
- The moduli space of abelian varieties of dimension  $g$  is a space of dimension  $g(g+1)/2$ . We have more liberty to find optimal abelian varieties in function of the security parameters.
- Supersingular abelian varieties can have larger embedding degree than supersingular elliptic curves.
- Over a Jacobian, we can use twists even if they are not coming from twists of the underlying curve.
- If  $A$  is an abelian variety of dimension  $g$ ,  $A[\ell]$  is a  $(\mathbb{Z}/\ell\mathbb{Z})$ -module of dimension  $2g \Rightarrow$  the structure of pairings on abelian varieties is richer.

## Polarised abelian varieties over $\mathbb{C}$

### Definition

A complex abelian variety  $A$  of dimension  $g$  is isomorphic to a compact Lie group  $V/\Lambda$  with

- A complex vector space  $V$  of dimension  $g$ ;
- A  $\mathbb{Z}$ -lattice  $\Lambda$  in  $V$  (of rank  $2g$ );

such that there exists an Hermitian form  $H$  on  $V$  with  $E(\Lambda, \Lambda) \subset \mathbb{Z}$  where  $E = \text{Im } H$  is symplectic.

- Such an Hermitian form  $H$  is called a **polarisation** on  $A$ . Conversely, any symplectic form  $E$  on  $V$  such that  $E(\Lambda, \Lambda) \subset \mathbb{Z}$  and  $E(ix, iy) = E(x, y)$  for all  $x, y \in V$  gives a polarisation  $H$  with  $E = \text{Im } H$ .
- Over a symplectic basis of  $\Lambda$ ,  $E$  is of the form.

$$\begin{pmatrix} 0 & D_{\delta} \\ -D_{\delta} & 0 \end{pmatrix}$$

where  $D_{\delta}$  is a diagonal positive integer matrix  $\delta = (\delta_1, \delta_2, \dots, \delta_g)$ , with  $\delta_1 | \delta_2 | \dots | \delta_g$ .

- The product  $\prod \delta_i$  is the degree of the polarisation;  $H$  is a **principal polarisation** if this degree is 1.

## Projective embeddings

### Proposition

Let  $\Phi : A = V/\Lambda \mapsto \mathbb{P}^{m-1}$  be a projective embedding. Then the linear functions  $f$  associated to this embedding are  $\Lambda$ -automorphics:

$$f(x + \lambda) = a(\lambda, x)f(x) \quad x \in V, \lambda \in \Lambda;$$

for a fixed automorphy factor  $a$ :

$$a(\lambda + \lambda', x) = a(\lambda, x + \lambda')a(\lambda', x).$$

### Theorem (Appell-Humbert)

All automorphy factors are of the form

$$a(\lambda, x) = \pm e^{\pi(H(x, \lambda) + \frac{1}{2}H(\lambda, \lambda))}$$

for a polarisation  $H$  on  $A$ .

## Theta functions

- Let  $(A, H_0)$  be a principally polarised abelian variety over  $\mathbb{C}$ :  
 $A = \mathbb{C}^g / (\Omega\mathbb{Z}^g + \mathbb{Z}^g)$  with  $\Omega \in \mathfrak{H}_g$ .
- The associated Riemann form on  $A$  is then given by  
 $E_1(\Omega x_1 + x_2, \Omega y_1 + y_2) = {}^t x_1 \cdot y_2 - {}^t y_1 \cdot x_2$ ; equivalently the matrix of  $H_0$  is  $\text{Im}\Omega^{-1}$ .
- The Weil pairing on  $A[\ell]$  corresponds to the symplectic form  $E$  on  $\frac{1}{\ell}\Lambda/\Lambda$ .
- All automorphic forms corresponding to a multiple  $H = nH_0$  of  $H_0$  come from the theta functions with characteristics:

$$\vartheta \left[ \begin{smallmatrix} a \\ b \end{smallmatrix} \right] (z, \Omega) = \sum_{n \in \mathbb{Z}^g} e^{\pi i {}^t (n+a)\Omega(n+a) + 2\pi i {}^t (n+a)(z+b)} \quad a, b \in \mathbb{Q}^g$$

- Automorphic property:

$$\vartheta \left[ \begin{smallmatrix} a \\ b \end{smallmatrix} \right] (z + m_1\Omega + m_2, \Omega) = e^{2\pi i ({}^t a \cdot m_2 - {}^t b \cdot m_1) - \pi i {}^t m_1 \Omega m_1 - 2\pi i {}^t m_1 \cdot z} \vartheta \left[ \begin{smallmatrix} a \\ b \end{smallmatrix} \right] (z, \Omega).$$

### Remark

*Working on level  $n$  mean we take a  $n$ -th power of the principal polarization. So in the following we will compute the  $n$ -th power of the usual Weil and Tate pairings.*

## Theta functions of level $n$

- Define  $\vartheta_i = \vartheta\left[\begin{smallmatrix} 0 \\ i \end{smallmatrix}\right]\left(\cdot, \frac{\Omega}{n}\right)$  for  $i \in Z(\bar{n}) = \mathbb{Z}^g/n\mathbb{Z}^g$  and
- This is a basis of the automorphic functions for  $H = nH_0$  (theta functions of level  $n$ );
- This is the unique basis such that in the projective coordinates:

$$\begin{aligned} A &\longrightarrow \mathbb{P}_{\mathbb{C}}^{n^g-1} \\ z &\longmapsto (\vartheta_i(z))_{i \in Z(\bar{n})} \end{aligned}$$

the translation by a point of  $n$ -torsion is normalized by

$$\vartheta_i\left(z + \frac{m_1}{n}\Omega + \frac{m_2}{n}\right) = e^{-\frac{2\pi i}{n} t_i \cdot m_1} \vartheta_{i+m_2}(z).$$

- $(\vartheta_i)_{i \in Z(\bar{n})} = \begin{cases} \text{coordinates system} & n \geq 3 \\ \text{coordinates on the Kummer variety } A/\pm 1 & n = 2 \end{cases}$
- $(\vartheta_i)_{i \in Z(\bar{n})}$ : basis of the theta functions of level  $n$   
 $\Leftrightarrow A[n] = A_1[n] \oplus A_2[n]$ : symplectic decomposition.
- Theta null point:  $\vartheta_i(0)_{i \in Z(\bar{n})} = \text{modular invariant}$ .

# Jacobians

- Let  $C$  be a curve of genus  $g$ ;
- Let  $V$  be the dual of the space  $V^* = \Omega^1(C, \mathbb{C})$  of holomorphic differentials of the first kind on  $C$ ;
- Let  $\Lambda \simeq H^1(C, \mathbb{Z}) \subset V$  be the set of periods (integration of differentials on loops);
- The intersection pairing gives a symplectic form  $E$  on  $\Lambda$ ;
- Let  $H$  be the associated hermitian form on  $V$ ;

$$H^*(w_1, w_2) = \int_C w_1 \wedge w_2;$$

- Then  $(V/\Lambda, H)$  is a principally polarised abelian variety: the **Jacobian** of  $C$ .

## Theorem (Torelli)

Jac  $C$  with the associated *principal polarisation* uniquely determines  $C$ .

## Remark (Weil pairing)

*In this setting, the Weil pairing can be seen as the intersection pairing on*

$$\text{Jac } C[\ell] \simeq \frac{1}{\ell} H_1(C, \mathbb{Z}) / H_1(C, \mathbb{Z}) \simeq H_1(C, \mathbb{Z}/\ell\mathbb{Z}).$$

# The differential addition law ( $k = \mathbb{C}$ )

$$\left( \sum_{t \in Z(\bar{2})} \chi(t) \vartheta_{i+t}(\mathbf{x} + \mathbf{y}) \vartheta_{j+t}(\mathbf{x} - \mathbf{y}) \right) \cdot \left( \sum_{t \in Z(\bar{2})} \chi(t) \vartheta_{k+t}(\mathbf{0}) \vartheta_{l+t}(\mathbf{0}) \right) =$$

$$\left( \sum_{t \in Z(\bar{2})} \chi(t) \vartheta_{-i'+t}(\mathbf{y}) \vartheta_{j'+t}(\mathbf{y}) \right) \cdot \left( \sum_{t \in Z(\bar{2})} \chi(t) \vartheta_{k'+t}(\mathbf{x}) \vartheta_{l'+t}(\mathbf{x}) \right).$$

where  $\chi \in \hat{Z}(\bar{2}), i, j, k, l \in Z(\bar{n})$

$$(i', j', k', l') = A(i, j, k, l)$$

$$A = \frac{1}{2} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & 1 & -1 \\ 1 & -1 & -1 & 1 \end{pmatrix}$$

## Example: differential addition in dimension 1 and in level 2

### Algorithm

**Input**  $z_P = (x_0, x_1)$ ,  $z_Q = (y_0, y_1)$  and  $z_{P-Q} = (z_0, z_1)$  with  $z_0 z_1 \neq 0$ ;  
 $z_0 = (a, b)$  and  $A = 2(a^2 + b^2)$ ,  $B = 2(a^2 - b^2)$ .

**Output**  $z_{P+Q} = (t_0, t_1)$ .

$$\textcircled{1} \quad t'_0 = (x_0^2 + x_1^2)(y_0^2 + y_1^2)/A$$

$$\textcircled{2} \quad t'_1 = (x_0^2 - x_1^2)(y_0^2 - y_1^2)/B$$

$$\textcircled{3} \quad t_0 = (t'_0 + t'_1)/z_0$$

$$\textcircled{4} \quad t_1 = (t'_0 - t'_1)/z_1$$

**Return**  $(t_0, t_1)$

## Miller functions with theta coordinates

### Proposition (Lubicz-R. [LR14])

- For  $P \in A$  we note  $z_p$  a lift to  $\mathbb{C}^g$ . We call  $P$  a projective point and  $z_p$  an affine point (because we describe them via their projective, resp affine, theta coordinates);
- We have (up to a constant)

$$f_{\lambda,P}(z) = \frac{\vartheta(z)}{\vartheta(z + \lambda z_p)} \left( \frac{\vartheta(z + z_p)}{\vartheta(z)} \right)^\lambda;$$

- So (up to a constant)

$$f_{\lambda,\mu,P}(z) = \frac{\vartheta(z + \lambda z_p) \vartheta(z + \mu z_p)}{\vartheta(z) \vartheta(z + (\lambda + \mu) z_p)}.$$

## Three way addition

### Proposition (Lubicz-R. [LR14])

From the affine points  $z_P, z_Q, z_R, z_{P+Q}, z_{P+R}$  and  $z_{Q+R}$  one can compute the affine point  $z_{P+Q+R}$ .

### Proof.

We can compute the three way addition using a generalised version of Riemann's relations:

$$\left( \sum_{t \in \mathbb{Z}(\bar{2})} \chi(t) \vartheta_{i+t}(z_{P+Q+R}) \vartheta_{j+t}(z_P) \right) \cdot \left( \sum_{t \in \mathbb{Z}(\bar{2})} \chi(t) \vartheta_{k+t}(z_Q) \vartheta_{l+t}(z_R) \right) =$$

$$\left( \sum_{t \in \mathbb{Z}(\bar{2})} \chi(t) \vartheta_{-i'+t}(z_0) \vartheta_{j'+t}(z_{Q+R}) \right) \cdot \left( \sum_{t \in \mathbb{Z}(\bar{2})} \chi(t) \vartheta_{k'+t}(z_{P+R}) \vartheta_{l'+t}(z_{P+Q}) \right).$$



## Three way addition in dimension 1 level 2

### Algorithm

Input *The points*  $x, y, z, X = y + z, Y = x + z, Z = x + y$ ;

Output  $T = x + y + z$ .

Return

$$T_0 = \frac{(aX_0 + bX_1)(Y_0Z_0 + Y_1Z_1)}{x_0(y_0z_0 + y_1z_1)} + \frac{(aX_0 - bX_1)(Y_0Z_0 - Y_1Z_1)}{x_0(y_0z_0 - y_1z_1)}$$

$$T_1 = \frac{(aX_0 + bX_1)(Y_0Z_0 + Y_1Z_1)}{x_1(y_0z_0 + y_1z_1)} - \frac{(aX_0 - bX_1)(Y_0Z_0 - Y_1Z_1)}{x_1(y_0z_0 - y_1z_1)}$$

## Computing the Miller function $f_{\lambda,\mu,P}((Q) - (0))$

### Algorithm

Input  $\lambda P$ ,  $\mu P$  and  $Q$ ;

Output  $f_{\lambda,\mu,P}((Q) - (0))$

- 1 Compute  $(\lambda + \mu)P$ ,  $Q + \lambda P$ ,  $Q + \mu P$  using normal additions and take any affine lifts  $z_{(\lambda+\mu)P}$ ,  $z_{Q+\lambda P}$  and  $z_{Q+\mu P}$ ;
- 2 Use a three way addition to compute  $z_{Q+(\lambda+\mu)P}$ ;

Return

$$f_{\lambda,\mu,P}((Q) - (0)) = \frac{\vartheta(z_Q + \lambda z_P)\vartheta(z_Q + \mu z_P)}{\vartheta(z_Q)\vartheta(z_Q + (\lambda + \mu)z_P)} \cdot \frac{\vartheta((\lambda + \mu)z_P)\vartheta(z_P)}{\vartheta(\lambda z_P)\vartheta(\mu z_P)}.$$

### Lemma

The result does not depend on the choice of affine lifts in Step 2.

- ☺ This allows us to evaluate the Weil and Tate pairings and derived pairings;
- ☹ Not possible *a priori* to apply this algorithm in level 2.

## The Tate pairing with Miller's functions and theta coordinates

- Let  $P \in A[\ell](\mathbb{F}_{q^d})$  and  $Q \in A(\mathbb{F}_{q^d})$ ; choose any lift  $z_P, z_Q$  and  $z_{P+Q}$ .
- The algorithm loop over the binary expansion of  $\ell$ , and at each step does a doubling step, and if necessary an addition step.

Given  $z_{\lambda P}, z_{\lambda P+Q}$ ;

Doubling Compute  $z_{2\lambda P}, z_{2\lambda P+Q}$  using two differential additions;

Addition Compute  $(2\lambda + 1)P$  and take an arbitrary lift  $z_{(2\lambda+1)P}$ . Use a three way addition to compute  $z_{(2\lambda+1)P+Q}$ .

- At the end we have computed affine points  $z_{\ell P}$  and  $z_{\ell P+Q}$ . Evaluating the Miller function then gives exactly the quotient of the projective factors between  $z_{\ell P}, z_0$  and  $z_{\ell P+Q}, z_Q$ .
- ☺ Described this way can be extended to level 2 by using **compatible additions**;
- ☹ Three way additions and normal (or compatible) additions are quite cumbersome, is there a way to only use differential additions?

## The Weil and Tate pairing with theta coordinates (Lubicz-R. [LR10])

Using directly the formula for  $f_{\ell,P}(z)$  we get that the Weil and Tate pairings are given by

$$e_{W,\ell}(P, Q) = \frac{\vartheta(z_Q + \ell z_P) \vartheta(0)}{\vartheta(z_Q) \vartheta(\ell z_P)} \frac{\vartheta(z_P) \vartheta(\ell z_Q)}{\vartheta(z_P + \ell z_Q) \vartheta(0)}$$

$$e_{T,\ell}(P, Q) = \frac{\vartheta(z_Q + \ell z_P) \vartheta(0)}{\vartheta(z_Q) \vartheta(\ell z_P)}$$

# The Weil and Tate pairing with theta coordinates (Lubicz-R. [LR10])

$P$  and  $Q$  points of  $\ell$ -torsion.

$z_0$	$z_P$	$2z_P$	$\dots$	$\ell z_P = \lambda_P^0 z_0$
$z_Q$	$z_P \oplus z_Q$	$2z_P + z_Q$	$\dots$	$\ell z_P + z_Q = \lambda_P^1 z_Q$
$2z_Q$	$z_P + 2z_Q$			
$\dots$	$\dots$			

$$\ell Q = \lambda_Q^0 0_A \quad z_P + \ell z_Q = \lambda_Q^1 z_P$$

- $e_{W,\ell}(P, Q) = \frac{\lambda_P^1 \lambda_Q^0}{\lambda_P^0 \lambda_Q^1}$ .
- $e_{T,\ell}(P, Q) = \frac{\lambda_P^1}{\lambda_P^0}$ .

## Why does it work?

$$\begin{array}{ccccccc}
 z_0 & \alpha z_P & \alpha^4(2z_P) & \dots & \alpha^{\ell^2}(lz_P) = \lambda'_P{}^0 z_0 \\
 \beta z_Q & \gamma(z_P \oplus z_Q) & \frac{\gamma^2 \alpha^2}{\beta}(2z_P + z_Q) & \dots & \frac{\gamma^\ell \alpha^{\ell(\ell-1)}}{\beta^{\ell-1}}(lz_P + z_Q) = \lambda'_P{}^1 \beta z_Q \\
 \beta^4(2z_Q) & \frac{\gamma^2 \beta^2}{\alpha}(z_P + 2z_Q) & & & \\
 \dots & \dots & & & \\
 \beta^{\ell^2}(lz_Q) = \lambda'_Q{}^0 z_0 & \frac{\gamma^\ell \beta^{\ell(\ell-1)}}{\alpha^{\ell-1}}(z_P + \ell z_Q) = \lambda'_Q{}^1 \alpha z_P & & & 
 \end{array}$$

We then have

$$\lambda'_P{}^0 = \alpha^{\ell^2} \lambda_P^0, \quad \lambda'_Q{}^0 = \beta^{\ell^2} \lambda_Q^0, \quad \lambda'_P{}^1 = \frac{\gamma^\ell \alpha^{\ell(\ell-1)}}{\beta^\ell} \lambda_P^1, \quad \lambda'_Q{}^1 = \frac{\gamma^\ell \beta^{\ell(\ell-1)}}{\alpha^\ell} \lambda_Q^1,$$

$$e'_{w,\ell}(P, Q) = \frac{\lambda'_P{}^1 \lambda'_Q{}^0}{\lambda'_P{}^0 \lambda'_Q{}^1} = \frac{\lambda_P^1 \lambda_Q^0}{\lambda_P^0 \lambda_Q^1} = e_{w,\ell}(P, Q),$$

$$e'_{T,\ell}(P, Q) = \frac{\lambda'_P{}^1}{\lambda'_P{}^0} = \frac{\gamma^\ell}{\alpha^\ell \beta^\ell} \frac{\lambda_P^1}{\lambda_P^0} = \frac{\gamma^\ell}{\alpha^\ell \beta^\ell} e_{T,\ell}(P, Q).$$

## Ate pairing

- Let  $P \in G_2 = A[\ell] \cap \text{Ker}(\pi_q - [q])$  and  $Q \in G_1 = A[\ell] \cap \text{Ker}(\pi_q - 1)$ ;  $\lambda \equiv q \pmod{\ell}$ .
- In projective coordinates, we have  $\pi_q^d(P + Q) = \lambda^d P + Q = P + Q$ ;
- Of course, in affine coordinates,  $\pi_q^d(z_{P+Q}) \neq \lambda^d z_P + z_Q$ .
- But if  $\pi_q(z_{P+Q}) = C * (\lambda z_P + z_Q)$ , then  $C$  is exactly the (non reduced) ate pairing (up to a renormalisation)!

### Algorithm (Computing the ate pairing)

Input  $P \in G_2, Q \in G_1$ ;

- 1 Compute  $z_Q + \lambda z_P, \lambda z_P$  using differential additions;
- 2 Find the projective factors  $C_1$  and  $C_0$  such that  $z_Q + \lambda z_P = C_1 * \pi(z_{P+Q})$  and  $\lambda z_P = C_0 * \pi(z_P)$  respectively;

Return  $(C_1/C_0)^{\frac{q^d-1}{\ell}}$ .

## Optimal ate pairing

- Let  $\lambda = m\ell = \sum c_i q^i$  be a multiple of  $\ell$  with small coefficients  $c_i$ . ( $\ell \nmid m$ )
- The pairing

$$a_\lambda: G_2 \times G_1 \longrightarrow \mu_\ell$$

$$(P, Q) \longmapsto \left( \prod_i f_{c_i, P}(Q)^{q^i} \prod_i f_{\sum_{j>i} c_j q^j, c_i q^i, P}(Q) \right)^{(q^d - 1)/\ell}$$

is non degenerate when  $mdq^{d-1} \not\equiv (q^d - 1)/r \sum_i i c_i q^{i-1} \pmod{\ell}$ .

- Since  $\varphi_d(q) = 0 \pmod{\ell}$  we look at powers  $q, q^2, \dots, q^{\varphi(d)-1}$ .
- We can expect to find  $\lambda$  such that  $c_i \approx \ell^{1/\varphi(d)}$ .

## Optimal ate pairing with theta functions

### Algorithm (Computing the optimal ate pairing)

Input  $\pi_q(P) = [q]P$ ,  $\pi_q(Q) = Q$ ,  $\lambda = m\ell = \sum c_i q^i$ ;

- 1 Compute the  $z_Q + c_i z_P$  and  $c_i q^i z_P$ ;
- 2 Apply Frobeniuses to obtain the  $z_Q + c_i q^i z_P$ ,  $c_i q^i z_P$ ;
- 3 Compute  $c_i q^i z_P \oplus \sum_j c_j q^j z_P$  (up to a constant) and then do a three way addition to compute  $z_Q + c_i q^i z_P + \sum_j c_j q^j z_P$  (up to the same constant);
- 4 Recurse until we get  $\lambda z_P = C_0 * z_P$  and  $z_Q + \lambda z_P = C_1 * z_Q$ ;

Return  $(C_1/C_0)^{\frac{q^d-1}{\ell}}$ .

## The case $n = 2$

- If  $n = 2$  we work over the Kummer variety  $K$  over  $k$ , so  $e(P, Q) \in \overline{k}^{\ast, \pm 1}$ .
- We represent a class  $x \in \overline{k}^{\ast, \pm 1}$  by  $x + 1/x \in \overline{k}^{\ast}$ . We want to compute the symmetric pairing

$$e_s(P, Q) = e(P, Q) + e(-P, Q).$$

- From  $\pm P$  and  $\pm Q$  we can compute  $\{\pm(P+Q), \pm(P-Q)\}$  (need a square root), and from these points the symmetric pairing.
- $e_s$  is compatible with the  $\mathbb{Z}$ -structure on  $K$  and  $\overline{k}^{\ast, \pm 1}$ .
- The  $\mathbb{Z}$ -structure on  $\overline{k}^{\ast, \pm 1}$  can be computed as follow:

$$\left(x^{\ell_1 + \ell_2} + \frac{1}{x^{\ell_1 + \ell_2}}\right) + \left(x^{\ell_1 - \ell_2} + \frac{1}{x^{\ell_1 - \ell_2}}\right) = \left(x^{\ell_1} + \frac{1}{x^{\ell_1}}\right) \left(x^{\ell_2} + \frac{1}{x^{\ell_2}}\right)$$

## Optimal pairings on Kummer varieties

- Computing  $c_i q^i z_p \pm \sum_j c_j q^j z_p$  requires a square root (very costly);
- And we need to recognize  $c_i q^i z_p + \sum_j c_j q^j z_p$  from  $c_i q^i z_p - \sum_j c_j q^j z_p$ .
- We will use **compatible additions**: if we know  $x, y, z$  and  $x+z, y+z$ , we can compute  $x+y$  without a square root;
- We apply the compatible additions with  $x = c_i q^i z_p, y = \sum_j c_j q^j z_p$  and  $z = z_Q$ .

## Compatible additions

- Recall that we know  $x, y, z$  and  $x + z, y + z$ ;
- From it we can compute  $(x + z) \pm (y + z) = \{x + y + 2z, x - y\}$  and of course  $\{x + y, x - y\}$ ;
- Then  $x + y$  is the element in  $\{x + y, x - y\}$  not appearing in the preceding set;
- Since  $x - y$  is a common point, we can recover it without computing a square root.

# The compatible addition algorithm in dimension 1

## Algorithm

Input  $x, y, Y = x + z, X = y + z;$

### 1 Computing $x \pm y$ :

$$\alpha = (x_0^2 + x_1^2)(y_0^2 + y_1^2)/A$$

$$\beta = (x_0^2 - x_1^2)(y_0^2 - y_1^2)/B$$

$$\kappa_{00} = (\alpha + \beta), \kappa_{11} = (\alpha - \beta)$$

$$\kappa_{10} := x_0 x_1 y_0 y_1 / ab$$

### 2 Computing $(x + z) \pm (y + z)$ :

$$\alpha' = (Y_0^2 + Y_1^2)(X_0^2 + X_1^2)/A$$

$$\beta' = (Y_0^2 - Y_1^2)(X_0^2 - X_1^2)/B$$

$$\kappa'_{00} = \alpha' + \beta', \kappa'_{11} = \alpha' - \beta'$$

$$\kappa'_{10} = Y_1 Y_2 X_1 X_2 / ab$$

Return  $x + y = [\kappa_{00}(\kappa_{10}\kappa'_{00} - \kappa'_{10}\kappa_{00}), \kappa_{10}(\kappa_{10}\kappa'_{00} - \kappa'_{10}\kappa_{00}) + \kappa_{00}(\kappa_{11}\kappa'_{00} - \kappa'_{11}\kappa_{00})]$

## One step of the pairing computation

### Algorithm (A step of the Miller loop with differential additions)

**Input**  $nP = (x_n, z_n)$ ;  $(n+1)P = (x_{n+1}, z_{n+1})$ ,  $(n+1)P + Q = (x'_{n+1}, z'_{n+1})$ .

**Output**  $2nP = (x_{2n}, z_{2n})$ ;  $(2n+1)P = (x_{2n+1}, z_{2n+1})$ ;  
 $(2n+1)P + Q = (x'_{2n+1}, z'_{2n+1})$ .

$$\textcircled{1} \quad \alpha = (x_n^2 + z_n^2); \beta = \frac{A}{B}(x_n^2 - z_n^2).$$

$$\textcircled{2} \quad X_n = \alpha^2; X_{n+1} = \alpha(x_{n+1}^2 + z_{n+1}^2); X'_{n+1} = \alpha(x'^2_{n+1} + z'^2_{n+1});$$

$$\textcircled{3} \quad Z_n = \beta(x_n^2 - z_n^2); Z_{n+1} = \beta(x_{n+1}^2 - z_{n+1}^2); Z'_{n+1} = \beta(x'^2_{n+1} + z'^2_{n+1});$$

$$\textcircled{4} \quad x_{2n} = X_n + Z_n; x_{2n+1} = (X_{n+1} + Z_{n+1})/x_P; x'_{2n+1} = (X'_{n+1} + Z'_{n+1})/x_Q;$$

$$\textcircled{5} \quad z_{2n} = \frac{a}{b}(X_n - Z_n); z_{2n+1} = (X_{n+1} - Z_{n+1})/z_P; z'_{2n+1} = (X'_{n+1} - Z'_{n+1})/z_Q;$$

**Return**  $(x_{2n}, z_{2n})$ ;  $(x_{2n+1}, z_{2n+1})$ ;  $(x'_{2n+1}, z'_{2n+1})$ .

## Weil and Tate pairing over $\mathbb{F}_{q^d}$

$g = 1$	$4\mathbf{M} + 2\mathbf{m} + 8\mathbf{S} + 3\mathbf{m}_0$
$g = 2$	$8\mathbf{M} + 6\mathbf{m} + 16\mathbf{S} + 9\mathbf{m}_0$

Table: Tate pairing with theta coordinates,  $P, Q \in A[\ell](\mathbb{F}_{q^d})$  (one step)

Operations in  $\mathbb{F}_q$ :  $\mathbf{M}$ : multiplication,  $\mathbf{S}$ : square,  $\mathbf{m}$  multiplication by a coordinate of  $P$  or  $Q$ ,  $\mathbf{m}_0$  multiplication by a theta constant;

Mixed operations in  $\mathbb{F}_q$  and  $\mathbb{F}_{q^d}$ :  $\mathbf{M}$ ,  $\mathbf{m}$  and  $\mathbf{m}_0$ ;

Operations in  $\mathbb{F}_{q^d}$ :  $\mathbf{M}$ ,  $\mathbf{m}$  and  $\mathbf{S}$ .

### Remark

- Doubling step for a Miller loop with Edwards coordinates:  $9\mathbf{M} + 7\mathbf{S} + 2\mathbf{m}_0$ ;
- Just doubling a point in Mumford projective coordinates using the fastest algorithm [HC]:  $21\mathbf{M} + 12\mathbf{S} + 2\mathbf{m}_0$ .
- Asymptotically the final exponentiation is more expensive than Miller's loop, so the Weil's pairing is faster than the Tate's pairing!

# Tate pairing

$g = 1$	$1\mathbf{m} + 2\mathbf{S} + 2\mathbf{M} + 2\mathbf{M} + 1\mathbf{m} + 6\mathbf{S} + 3\mathbf{m}_0$
$g = 2$	$3\mathbf{m} + 4\mathbf{S} + 4\mathbf{M} + 4\mathbf{M} + 3\mathbf{m} + 12\mathbf{S} + 9\mathbf{m}_0$

Table: Tate pairing with theta coordinates,  $P \in A[\ell](\mathbb{F}_q), Q \in A[\ell](\mathbb{F}_{q^d})$  (one step)

		Miller		Theta coordinates
		Doubling	Addition	One step
$g = 1$	$d$ even	$1\mathbf{M} + 1\mathbf{S} + 1\mathbf{M}$	$1\mathbf{M} + 1\mathbf{M}$	$1\mathbf{M} + 2\mathbf{S} + 2\mathbf{M}$
	$d$ odd	$2\mathbf{M} + 2\mathbf{S} + 1\mathbf{M}$	$2\mathbf{M} + 1\mathbf{M}$	
$g = 2$	$Q$ degenerate +	$1\mathbf{M} + 1\mathbf{S} + 3\mathbf{M}$	$1\mathbf{M} + 3\mathbf{M}$	$3\mathbf{M} + 4\mathbf{S} + 4\mathbf{M}$
	$d$ even			
	General case	$2\mathbf{M} + 2\mathbf{S} + 18\mathbf{M}$	$2\mathbf{M} + 18\mathbf{M}$	

Table:  $P \in A[\ell](\mathbb{F}_q), Q \in A[\ell](\mathbb{F}_{q^d})$  (counting only operations in  $\mathbb{F}_{q^d}$ ).

## Ate and optimal ate pairings

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$$g = 1 \quad 4\mathbf{M} + 1\mathbf{m} + 8\mathbf{S} + 1\mathbf{m} + 3\mathbf{m}_0$$

$$g = 2 \quad 8\mathbf{M} + 3\mathbf{m} + 16\mathbf{S} + 3\mathbf{m} + 9\mathbf{m}_0$$


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Table: Ate pairing with theta coordinates,  $P \in G_2, Q \in G_1$  (one step)

### Remark

Using affine Mumford coordinates in dimension 2, the hyperelliptic ate pairing costs [GHO+07]:

Doubling  $1\mathbf{I} + 29\mathbf{M} + 9\mathbf{S} + 7\mathbf{M}$

Addition  $1\mathbf{I} + 29\mathbf{M} + 5\mathbf{S} + 7\mathbf{M}$

(where  $\mathbf{I}$  denotes the cost of an affine inversion in  $\mathbb{F}_{q^d}$ ).

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