

Courbes elliptiques — 4TMA902U

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Identity Based Encryption

Boneh and Franklin have proposed in 2001 an identity based encryption (IBE) scheme using pairing friendly elliptic curves. Let $e : G \times G \rightarrow G_t$ be a type I (symmetric) cryptographic pairing. Let ℓ be the prime order of G and G_t and P a generator of G .

Disclaimer : we will use a type I pairing for simplicity, but in practice this means that we will use a supersingular curve that must be defined over a large prime field for security. This is inefficient compared with using a type III pairing.

1 Alice has a public key $A := aP$ and a secret key $a \in \mathbf{Z}/\ell\mathbf{Z}$. Bob has a public key $B := bP$ and a secret key $b \in \mathbf{Z}/\ell\mathbf{Z}$. Let $m \in G_t$ be a plaintext. Show how Carl can create a ciphertext C of m that Alice and Bob can both decrypt (hint : get inspired by the Tripartite Diffie-Hellman key exchange).

2 Oscar intercepts C . What problem must he solve in order to decrypt C and retrieve m ?

3 From this protocol, we want to devise an identity based encryption (IBE) scheme. To simplify, suppose that the set of identities are the elements of G . Suppose Carl wants to encrypt $m \in G_t$ for Bob using B as the identity of Bob (and assuming that Bob does not know the secret b anymore). From the protocol of question 1 show how to build this IBE scheme (hint : view Alice has the public key generator that provides a secret key to Bob in order to decrypt).

4 Now suppose that identities can be arbitrary bit-strings, $ID \in \{0, 1\}^*$. How can we get a full IBE scheme from the scheme of the previous question (this should give you the Boneh and Franklin IBE scheme)?

5 Implement this scheme using a supersingular curve.

6 Show how to simplify the generic construction from an IBE scheme to a digital signature scheme in order to retrieve the BLS scheme of last week from the Boneh and Franklin IBE scheme.