

Factoring linear differential operators in positive characteristic

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Algebra of differential operators

Object of study

Differential operators in $\mathbb{F}_p(x)\langle\partial\rangle = \{a_n(x)\partial^n + \cdots + a_1(x)\partial + a_0(x)\}$.

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$$\partial f = f\partial + f'$$

Derivation:

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Running example:

$$L = (2x^6 + 2)\partial^6 + (x^6 + 2)\partial^3 + 2x^6 + 2x^3 + 2 \in \mathbb{F}_3(x)\langle\partial\rangle$$

State of the art

- M. van der Put. Modular methods for factoring differential operators. Unpublished manuscript, 1997.
- M. Giesbrecht, Y. Zhang, Factoring and decomposing Ore polynomials over $\mathbb{F}_p(t)$, ISSAC 2003.
- T. Cluzeau, factorisation of differential systems in characteristic p , ISSAC 2003.
- X. Caruso, J. Le Borgne. A new faster algorithm for factoring skew polynomials over finite fields. JSC 2017.
- J. Gomez-Torrecillas, F. J. Lobillo, G. Navarro, Computing the bound of an Ore polynomial. Applications to factorisation, JSC 2019.

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Problem: For reducible divisors of $Q(x^p, \partial^p)$ with $Q \in \mathbb{F}_p[u, v]$.

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Notation

$C := \mathbb{F}_p(x^p)$ is the field of constants of $\mathbb{F}_p(x)$.

$C[\partial^p]$ is the center of $\mathbb{F}_p(x)\langle\partial\rangle$.

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Tool 1: p -curvature and first reduction (van der Put, Cluzeau)

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Contribution

A full factorisation algorithm that extends to operators with coefficients in finite separable extensions of $\mathbb{F}_p(x)$.

A guideline: studying the submodules of $\mathbb{F}_p(x)\langle\partial\rangle/\mathbb{F}_p(x)\langle\partial\rangle L$

Notation

- $\mathcal{D}_L := \mathbb{F}_p(x)\langle\partial\rangle/\mathbb{F}_p(x)\langle\partial\rangle L$.
- $\mathcal{D}_L L'$ is the $\mathbb{F}_p(x)\langle\partial\rangle$ -submodule of \mathcal{D}_L generated by L' .

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Direct consequence of $\mathbb{F}_p(x)\langle\partial\rangle$ being a left principal ideal domain.

p -curvature

Definition

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Proposition

Its characteristic polynomial $\chi(\psi_p^L)$ belongs to $\mathbf{C}[\mathcal{Y}]$.

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Operator version of classical classification results of differential modules in positive characteristic (van der Put).

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- $\mathcal{D}_{N(\partial^p)}$ is a C_N -algebra ($Y \mapsto \partial^p$)
- \mathcal{D}_L is a left $\mathcal{D}_{N(\partial^p)}$ -module

Structure results

Proposition (van der Put)

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Corollary

$\mathcal{D}_{N(\partial^p)}$ is either a division algebra or isomorphic to $M_p(C_N)$.

Morita's equivalence

If $\mathcal{D}_{N(\partial^p)}$ is a division algebra then $N(\partial^p)$ has no nontrivial divisor and $L = N(\partial^p)$ is irreducible.

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- If L' is a divisor of $N(\partial^p)$ then $\dim \text{Mor}_N(\mathcal{D}_{L'}) = \text{ord}(L') / \deg(N)$.

Running example

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Let $(H_i)_{1 \leq i \leq p}$ be irreducible divisors of $N(\partial^p)$ such that $N(\partial^p) = \text{lclm}_{i=1}^p H_i$.

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 For at least one i ,

$$\text{lclm}(H_i, R) \cdot R^{-1}$$

is an irreducible divisor of L .

Algorithm

- Compute $\chi(\psi_p^L)$ and factor it $\chi(\psi_p^L) = \prod_{i=1}^n N_i^{\nu_i}(Y)$.
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- Compute $\text{lclm}(R'_j, H_{i,k}) \cdot (R'_j)^{-1}$

Extending the field of constants

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Recall: $\mathcal{D}_{N(\partial^p)} = \mathbb{F}_p(x)\langle\partial\rangle/N(\partial^p)$

$$\begin{array}{ccc}
 \mathbb{F}_p(x)\langle\partial\rangle & \hookrightarrow & K_N\langle\partial\rangle \\
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 \mathbb{F}_p(x)\langle\partial\rangle/N(\partial^p) & \xrightarrow{\varphi_N} & K_N\langle\partial\rangle/\partial^p - y_N
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“ p -Riccati” equation

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Lemma (P., 2022)

Let $f \in K_N$ be such that $f^{(p-1)} + f^p = y_N$. Then

$$\partial^p - y_N = \text{lclm}_{i=1}^p \left(\partial - f - \frac{i}{x} \right)$$

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- For each $N_i(\partial^p)$ compute $(H_{i,k})_{k \in \llbracket 1;p \rrbracket}$ such that $\text{lclm}_{k=1}^p H_{i,k} = N_i(\partial^p)$.
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 - Compute $f_i \in K_N$ such that $f_i^{(p-1)} + f_i^p = y_{N_i}$.
- Set $H_{i,k} = \text{gcd}(\varphi_{N_i}^{-1}(\partial - f_i - \frac{k}{x}), N_i(\partial^p))$.
- Compute $\text{lclm}(R'_j, H_{i,k}) \cdot (R'_j)^{-1}$

The case when $L \neq N(\partial^p)$ (after van der Put)

$$L_N := \text{gcd}(\varphi_N(L), \partial^p - y_N)$$

is a non trivial divisor of $\partial^p - y_N$.

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Fact

If $L_N = \partial^m + a_{m-1}\partial^{m-1} + \cdots + a_0$ then $-\frac{a_{m-1}}{m}$ is a solution of the p -Riccati equation.

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If $L_N = \partial^m + a_{m-1}\partial^{m-1} + \cdots + a_0$ then $-\frac{a_{m-1}}{m}$ is a solution of the p -Riccati equation.

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The space of solutions of p -Riccati is an affine space

Algorithm

- Compute $\chi(\psi_p^L)$ and factor it $\chi(\psi_p^L) = \prod_{i=1}^n N_i^{r_i}(Y)$.
- Compute $L = L'_1 \cdots L'_m$ such that each L'_j divides a $N_i(\partial^p)$.
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 - **Compute $f_i \in K_N$ such that $f_i^{(p-1)} + f_i^p = y_{N_i}$.**
- Set $H_{i,k} = \text{gcd}(\varphi_{N_i}^{-1}(\partial - f - \frac{i}{x}), N_i(\partial^p))$.
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Factoring $N(\partial^p)$: when $K_N = \mathbb{F}_p(\mathbf{x})$

Suppose that $K_N = \mathbb{F}_p(\mathbf{x})$, $y_N = g^p \in \mathbb{F}_p(\mathbf{x}^p)$
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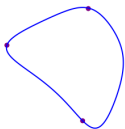
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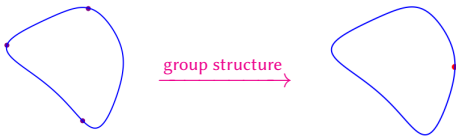
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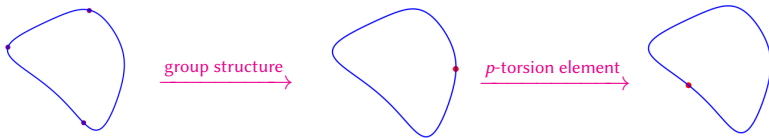
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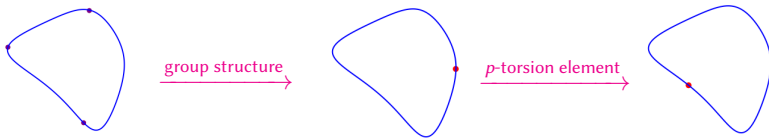
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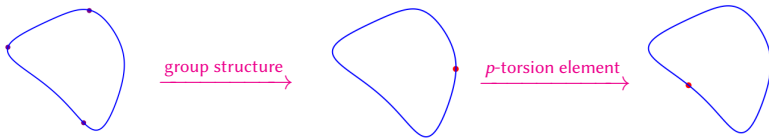
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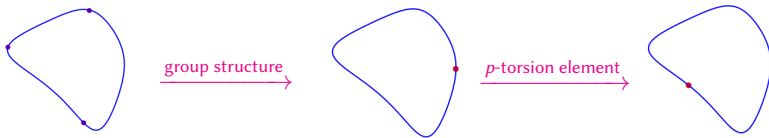
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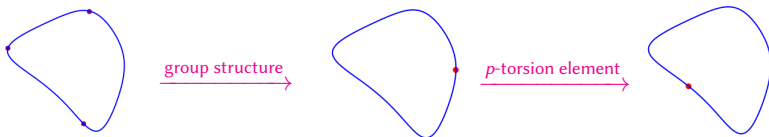
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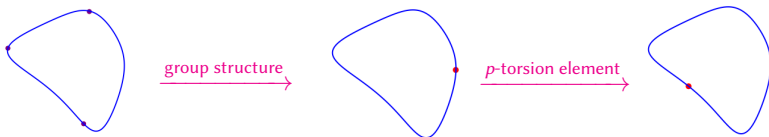
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Thank you for your attention